

Network Model

- OSI model and TCP/IP model -



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■ 학습개요

- 네트워크 모델에 대하여 이해하고, OSI-7 계층 모델과 TCI/IP 계층 모델의 구성을 학습한다.

■ 학습목표

- 네트워크 모델의 필요성과 개념을 설명할 수 있어야 한다.
- 네트워크 계층 모델에서 각 계층의 역할과 동작을 설명할 수 있어야 한다.

1. Layered Tasks

*We use the concept of **layers** in our daily life. As an example, let us consider two friends who communicate through postal mail. The process of sending a letter to a friend would be complex if there were no services available from the post office.*

Topics discussed in this section:

Sender, Receiver, and Carrier
Hierarchy

1. Layered Tasks

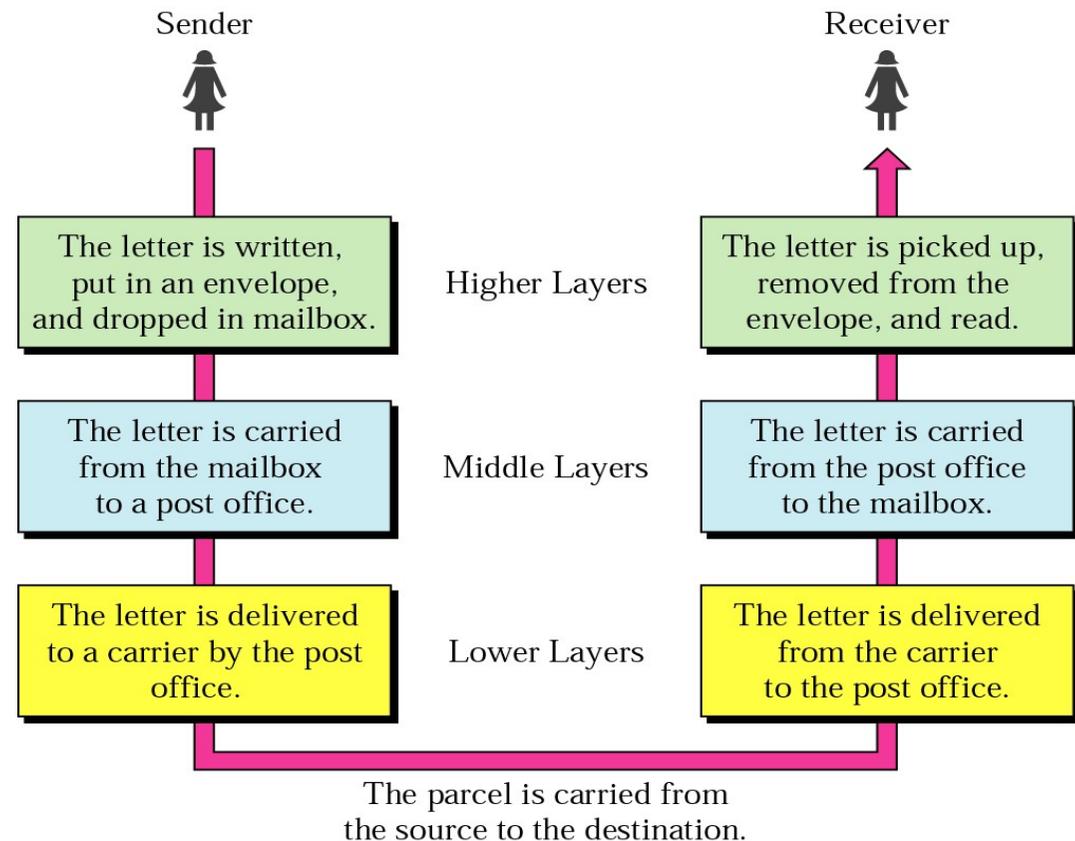
- We use the concept of layers in our daily life
 - ex : postal mail
- Sender, Receiver, and Carrier
 - needs three components
 - there is a hierarchy of tasks

- Hierarchy

- letter must be written and dropped in the mailbox before being picked up by the letter carrier at the sender
- letter must be dropped in the recipient mailbox before being picked up by the recipient

- Services

- each layer at the sending site uses the services of the layer immediately below it



2. The OSI Model

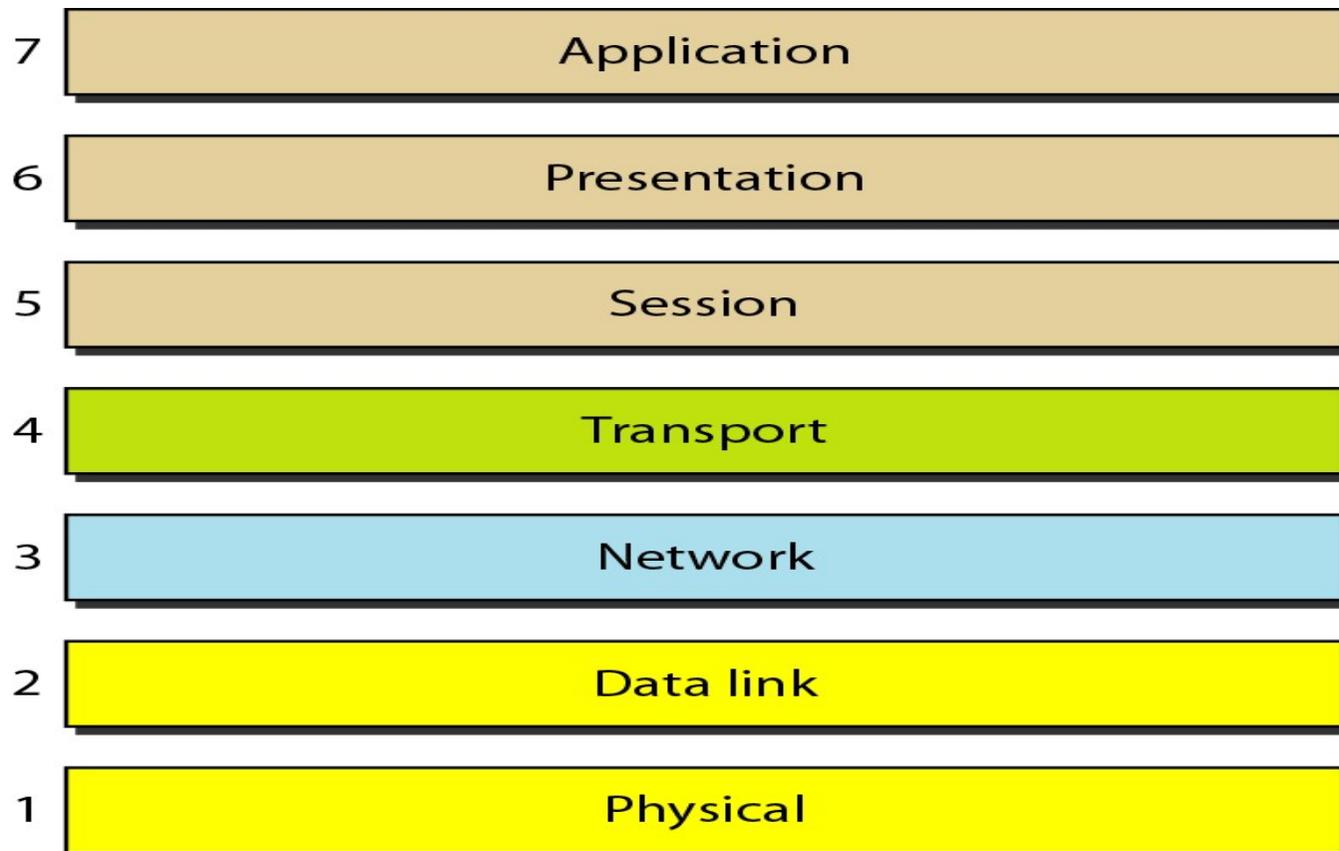
Established in 1947, the International Standards Organization (ISO) is a multinational body dedicated to worldwide agreement on international standards. An ISO standard that covers all aspects of network communications is the Open Systems Interconnection (OSI) model. It was first introduced in the late 1970s.

Topics discussed in this section:

Layered Architecture
Peer-to-Peer Processes
Encapsulation

2. The OSI Model

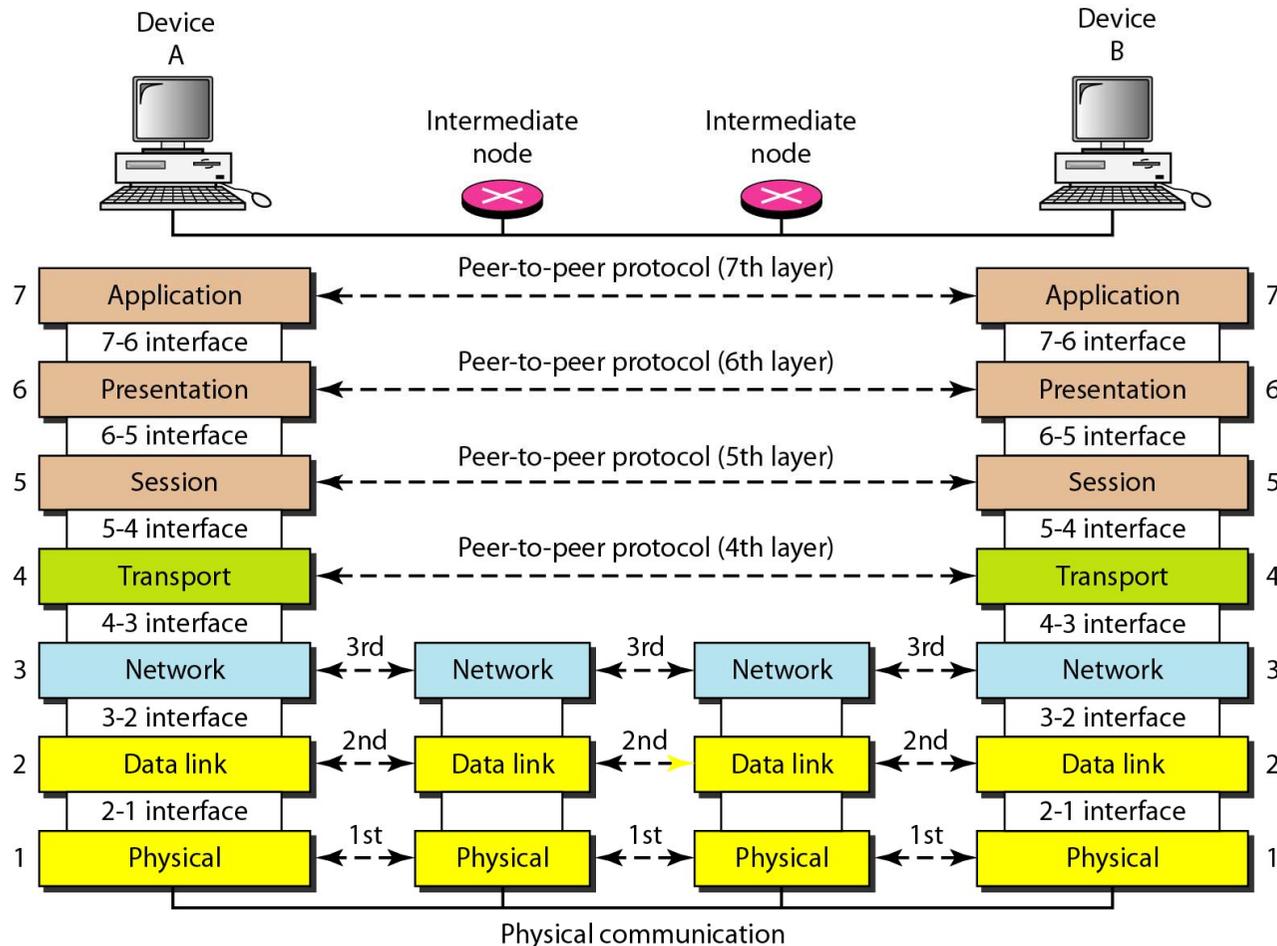
- The International Standards Organization (ISO) is a multinational body dedicated to worldwide agreement of international standards.
 - 7 layer
 - Layer defines a segment of the process of moving information across a network



2. The OSI Model

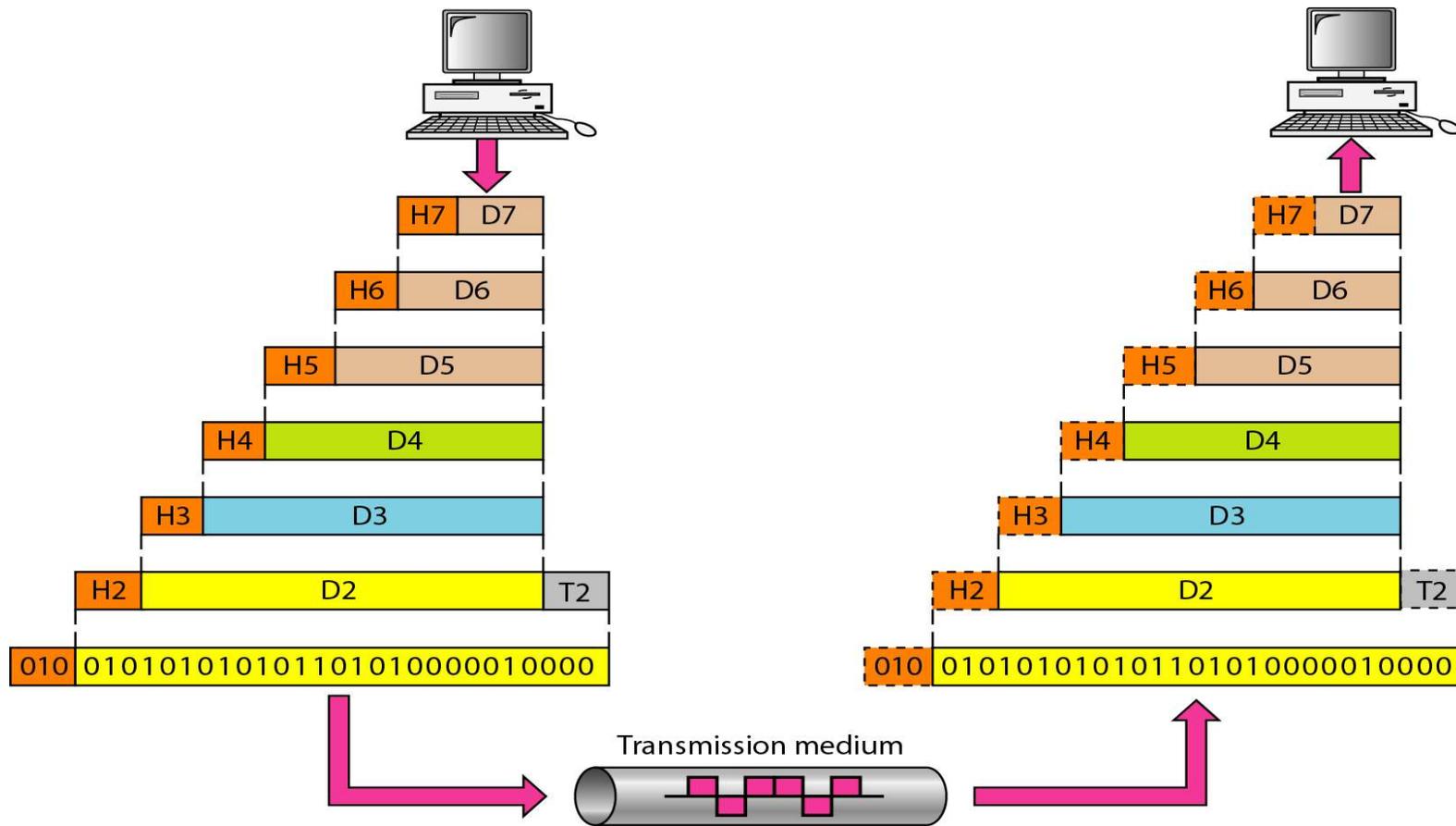
Network Models

- Layers involved when a message is sent from device A to device B
 - may pass through many intermediate nodes
 - intermediate nodes usually involve only the first three layers of the model



2. The OSI Model

Peer-to-Peer Process



2. The OSI Model

Network Models

- interfaces between layers
 - data passing is made possible by an interface between each pair of adjacent layers
 - defines what information and services a layer must provide for above layer
 - provide modularity → specific implementation of a layer's functions can be modified or replaced without requiring changes to the surrounding layers
- organization of the layers
 - L1, 2, and 3 → network support layer
 - deal with the physical aspects of moving data from device to another
 - such as electrical specifications, physical connections, physical addressing, transport timing and reliability
 - L5, 6, and 7 → user support layer
 - allows interoperability among unrelated software system
 - L4 → links the two subgroups
 - ensures end-to-end reliable data transmission
 - header can be added at each layer
 - trailer is added at layer 2
 - at the receiver, the headers(trailers) attached at the corresponding layer are removed, and action appropriate to that layer are taken

3. Layers In the OSI Models

In this section we briefly describe the functions of each layer in the OSI model.

Topics discussed in this section:

Physical Layer

Data Link Layer

Network Layer

Transport Layer

Session Layer

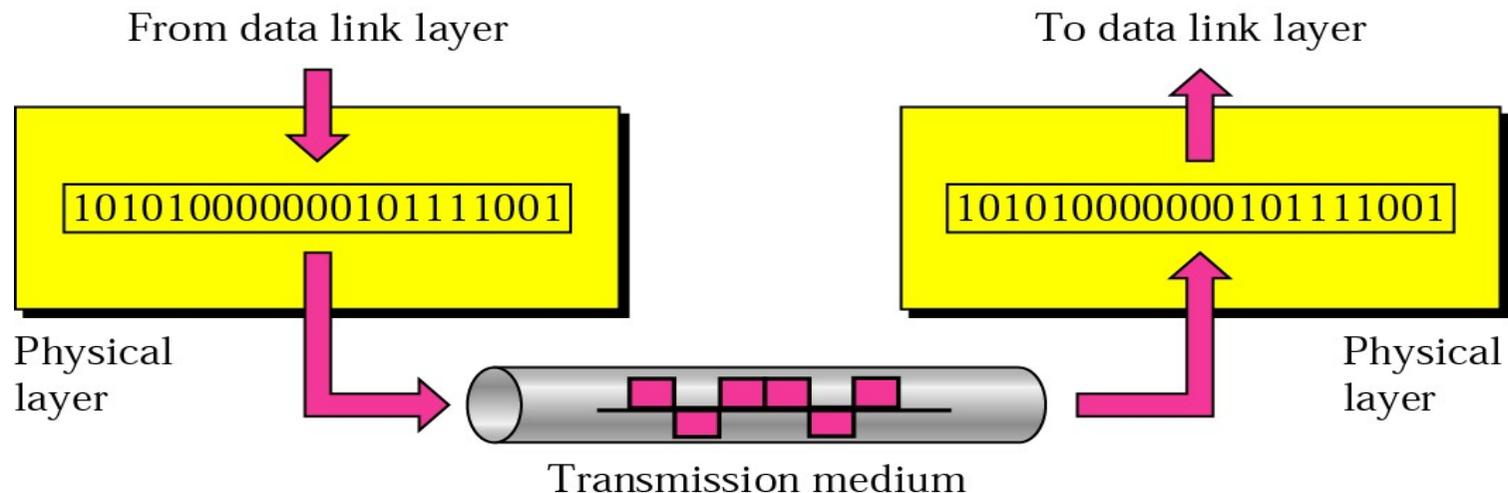
Presentation Layer

Application Layer

3. Layers In the OSI Models

1) Physical Layer

- coordinates the functions required to transmit a bit stream over a physical medium



- Characteristics of interfaces and medium
 - physical, electrical, procedural, and functional characteristics of the interface between the devices and the transmission medium
 - defines the type of the transmission medium
- Representation of bits
 - defines the type of encoding
 - how 0s and 1s are changed to signals – electrical or optical

3. Layers In the OSI Models

1) Physical Layer

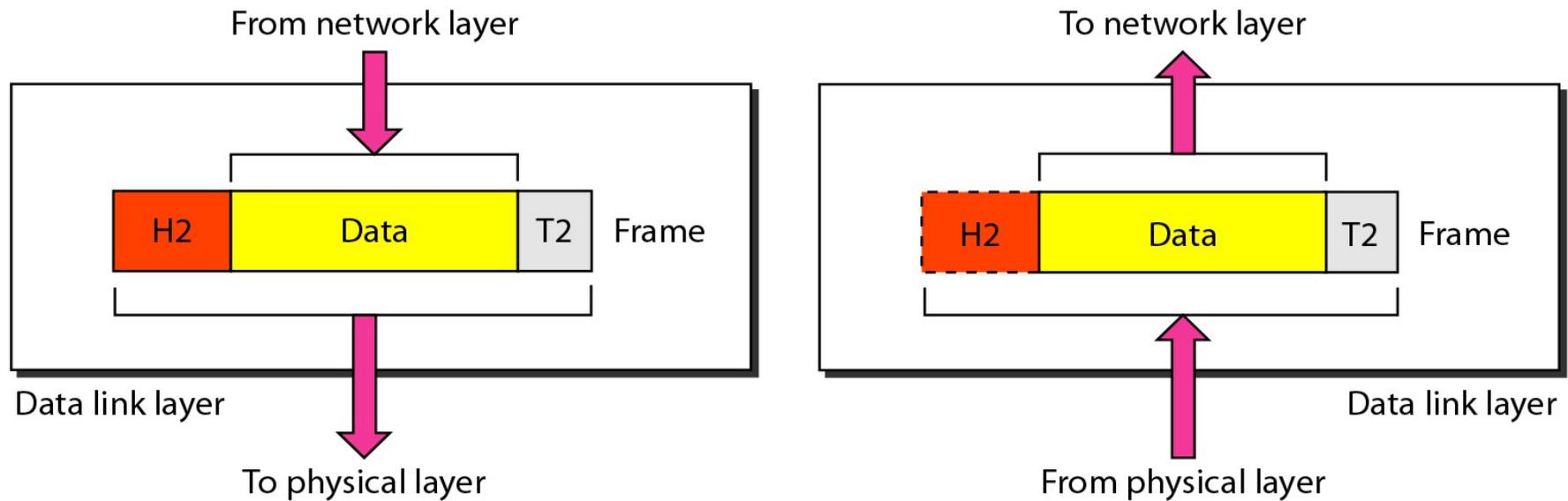
- Data rate, transmission rate : the number of bits sent each seconds
 - defines the duration of a bit
- Synchronization of bits
 - sender and receiver clocks must be synchronized
- Line configuration
 - The physical layer is concerned with the connection of device to the media
- Physical topology
 - The physical topology defines how devices are connected to make a network.
- Transmission mode
 - The physical layer also defines the direction of transmission between two devices: simplex, half-duplex, or full-duplex.

The physical layer is responsible for movements of individual bits from one hop (node) to the next.

3. Layers In the OSI Models

2) Data Link Layer

- makes the physical layer appear error-free to the upper (network) layer



3. Layers In the OSI Models

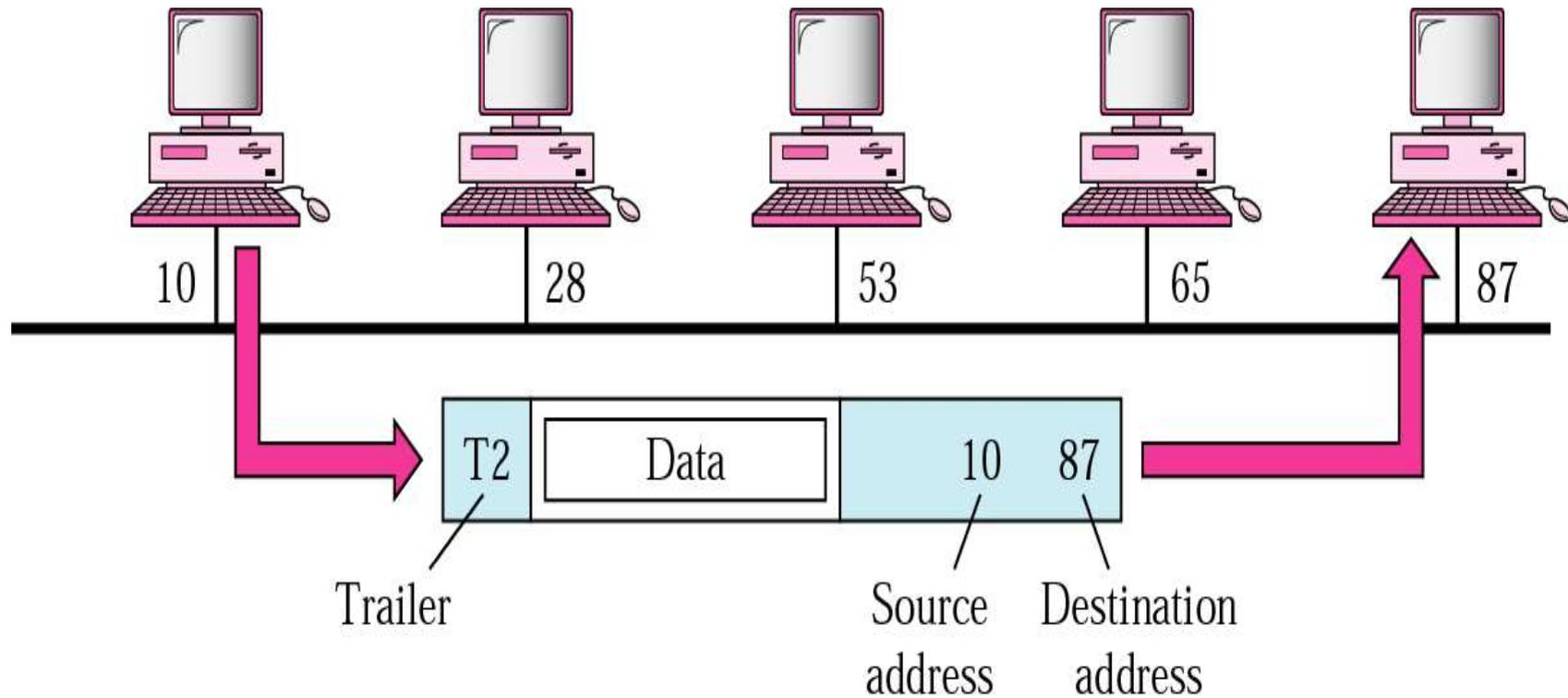
2) Data Link Layer

- responsible for node-to-node delivery
 - framing
 - divides the data stream into manageable data unit → frame
 - detection of frame at the receiver
 - physical addressing
 - adds a header to define the sender (source address) and/or receiver (destination address)
 - distribute frames to different systems on the sender's network
 - receiver address is the address of the devices that connects one network to the next if the frame is intended for a system outside the sender's network
 - flow control : necessary when data is being sent faster than it can be processed by receiver
 - error control
 - detect and retransmit damaged or lost frames to add reliability to the physical layer
 - also prevent duplication of frames
 - normally achieved through a trailer added to the end of the frames
 - access control : determines which device has control over the link at any given time when two or more devices are connected to the same link

3. Layers In the OSI Models

2) Data Link Layer

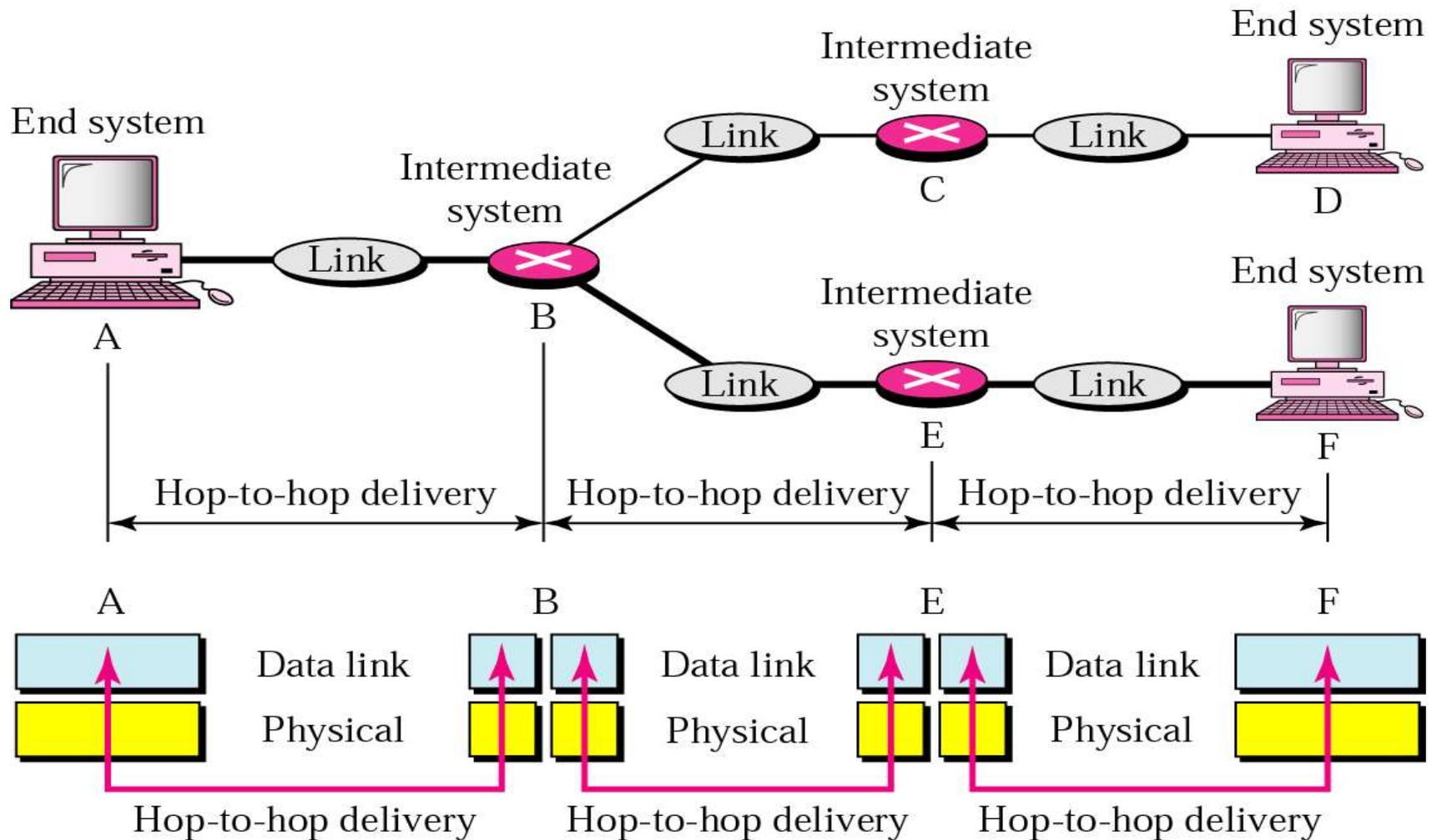
- Example.



3. Layers In the OSI Models

2) Data Link Layer

- hop-by-hop (node-to-node) delivery concept by the data link layer



3. Layers In the OSI Models

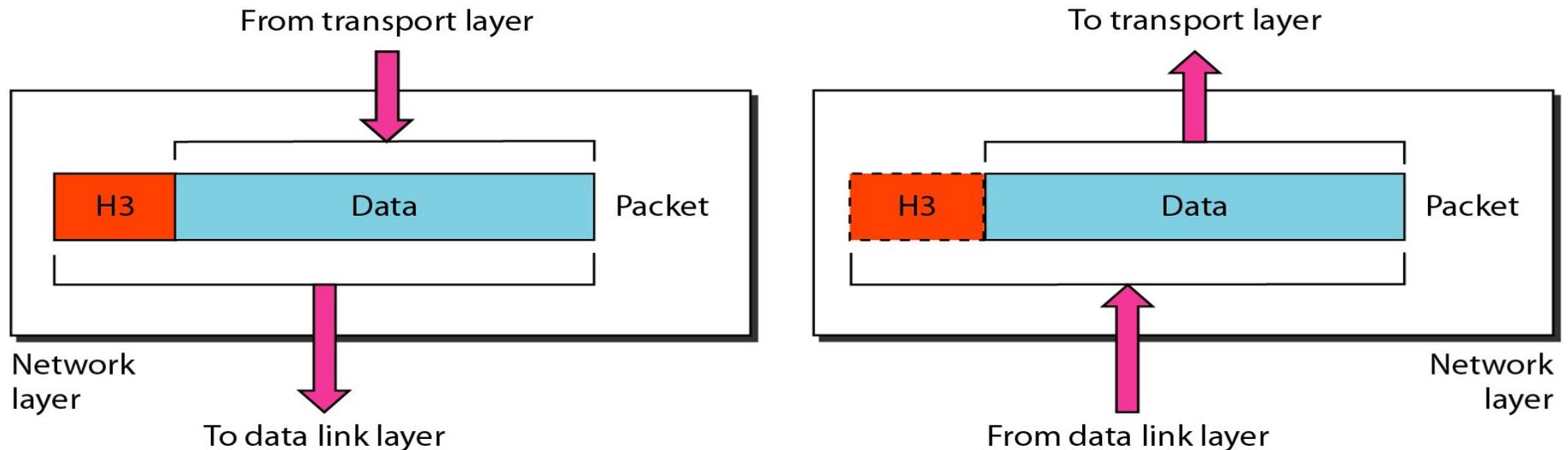
2) Data Link Layer

Note

The data link layer is responsible for moving frames from one hop (node) to the next.

3. Layers In the OSI Models

3) Network Layer

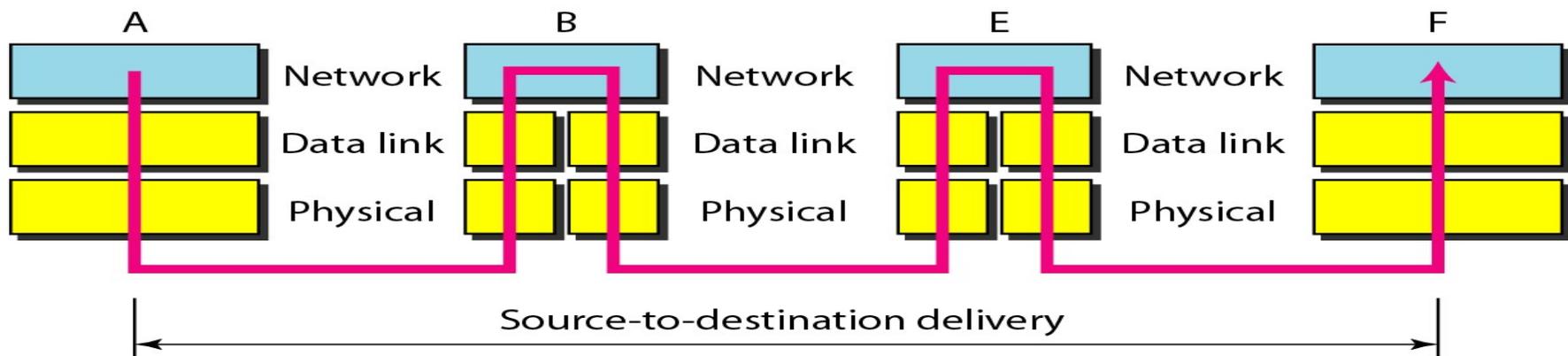
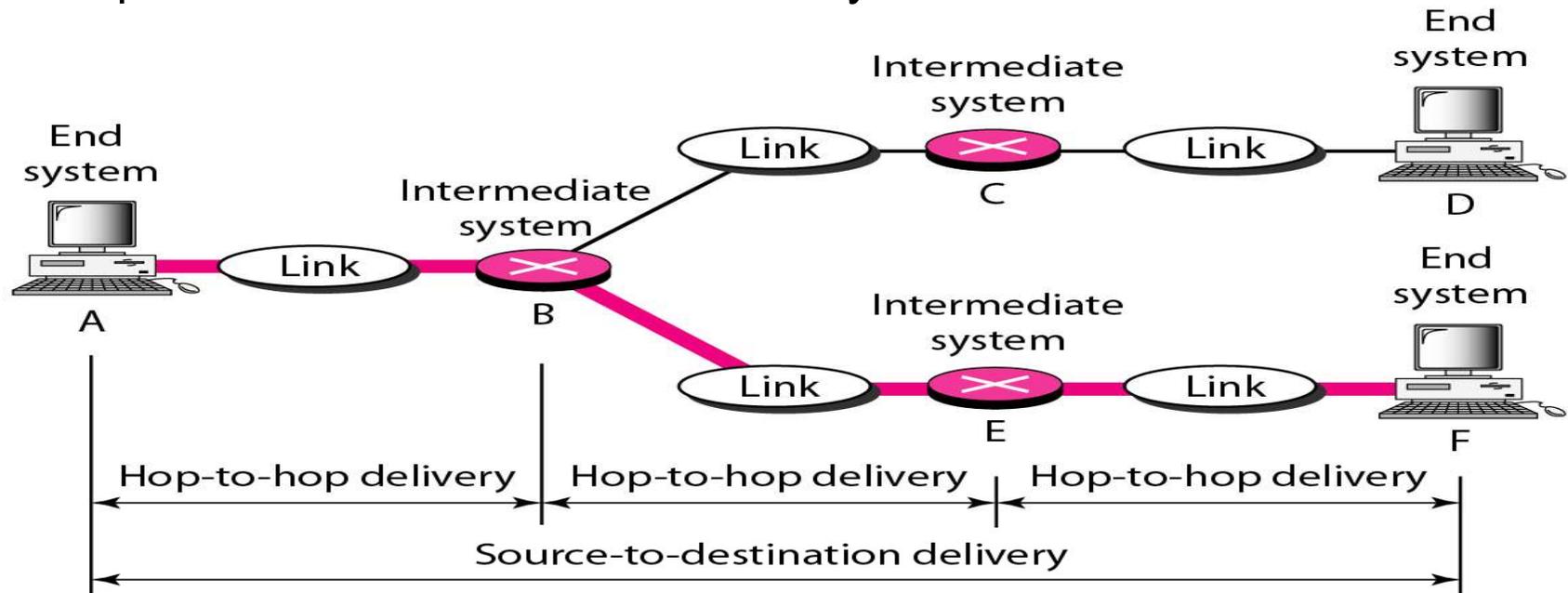


- responsible for the delivery of packets from the original source to the final destination possibly across multiple networks (links)
 - logical addressing
 - If a packet passes the network boundary, we need it to help distinguish the source and destination systems
 - routing
 - in the internet, connecting devices (called routers or switches) route or switch the packets to their final destination

3. Layers In the OSI Models

3) Network Layer

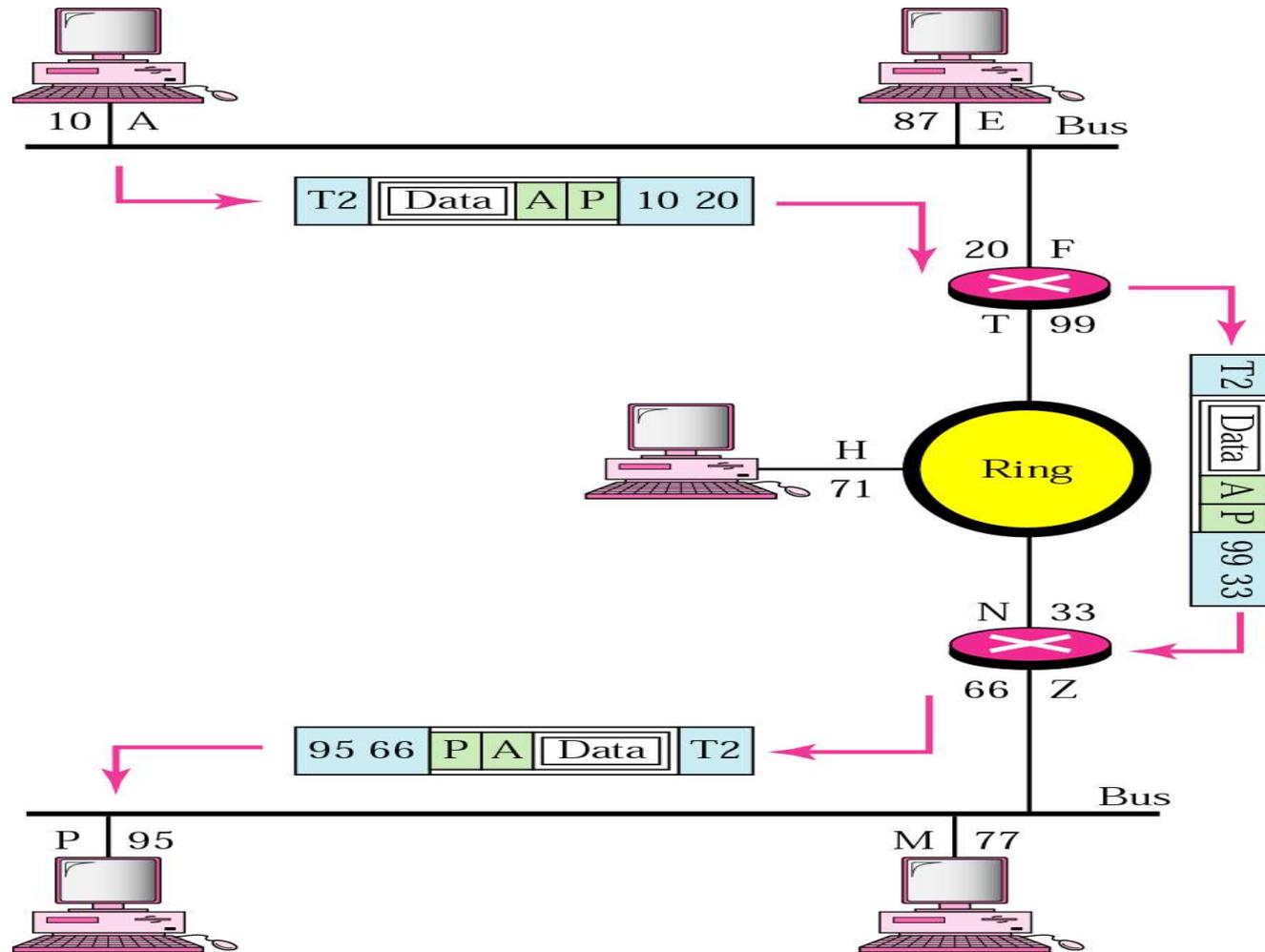
- concept of source-to-destination delivery



3. Layers In the OSI Models

3) Network Layer

■ Example.



3. Layers In the OSI Models

3) Network Layer

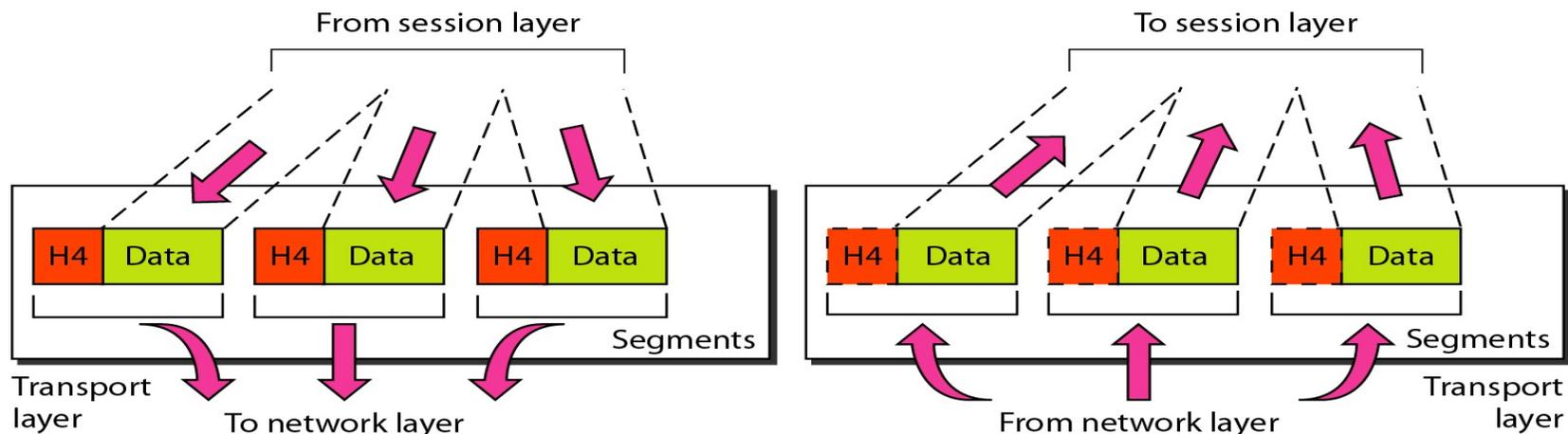


The network layer is responsible for the delivery of individual packets from the source host to the destination host.

3. Layers In the OSI Models

4) Transport Layer

- responsible for process-to-process delivery of the entire message
 - treats each packet independently
 - ensures that the whole message arrives intact and in order
 - use both error control and flow control at the source-to-destination level
- port addressing
 - computers often run several processes (running program) at the same time
 - process-to-process delivery means delivery to a specific process
 - header must therefore include a type of address called port address
- segmentation and reassembly
 - a message is divided into transmittable segments at the sender
 - reassemble the message correctly upon arriving at the receiver



3. Layers In the OSI Models

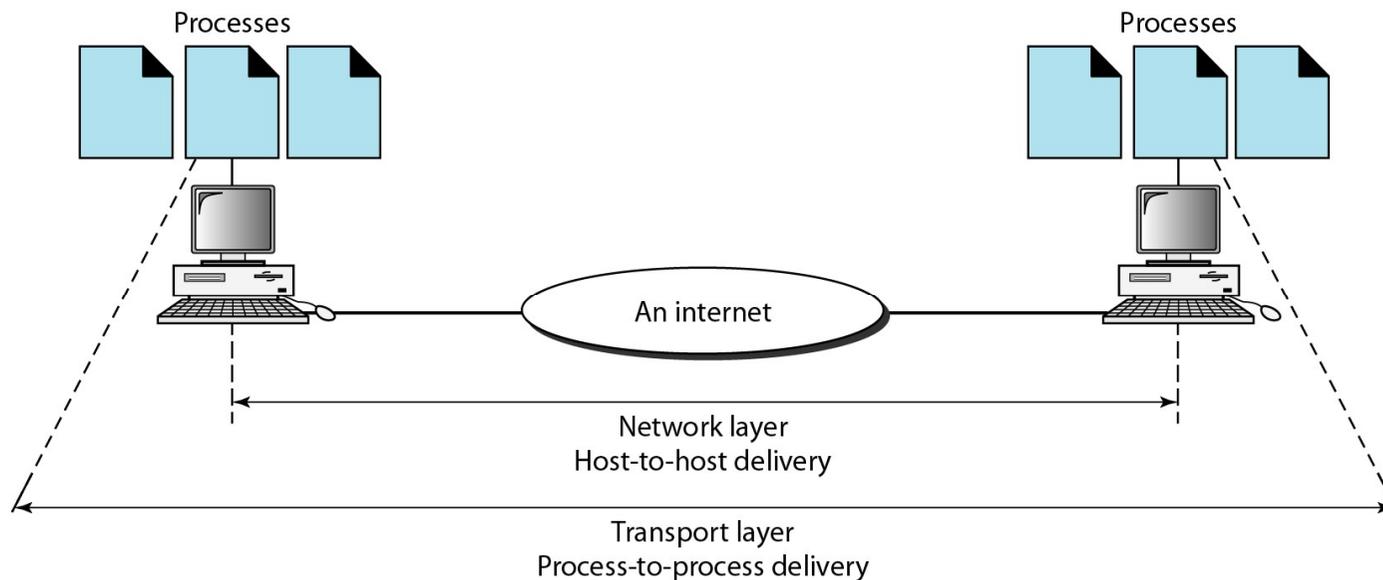
4) Transport Layer

■ connection control

- connectionless mode : treats each segment as an independent packet
- connection-oriented mode : needs connection control
 - make a connection with the transport layer at the destination machine before delivering the packets
 - after all the data are transferred, the connection is terminated
 - has more control over sequencing, flow and error control

■ flow control : end-to-end flow control rather than across a single link

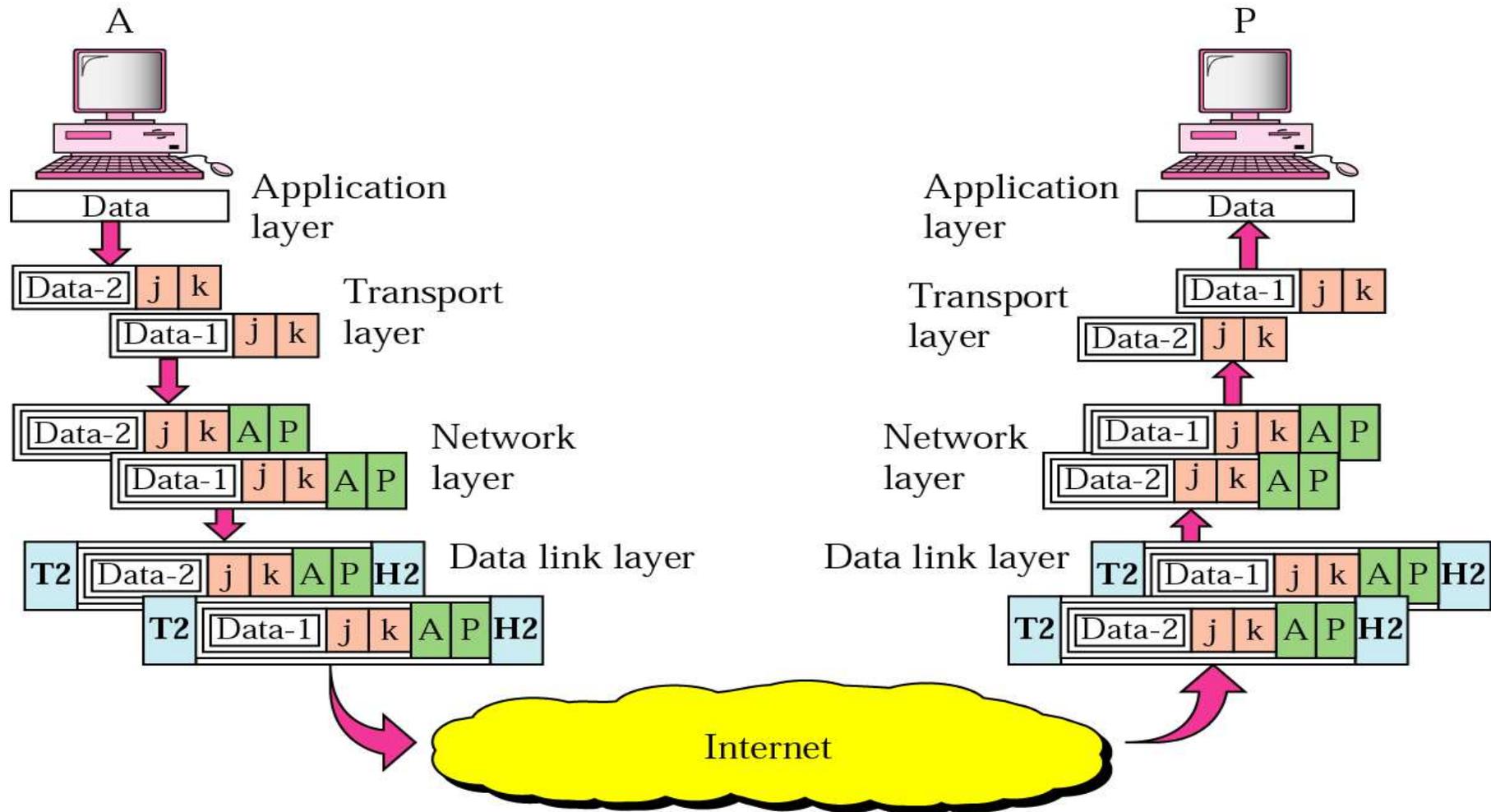
■ error control : end-to-end error control rather than across a single link



3. Layers In the OSI Models

4) Transport Layer

■ Example



3. Layers In the OSI Models

4) Transport Layer

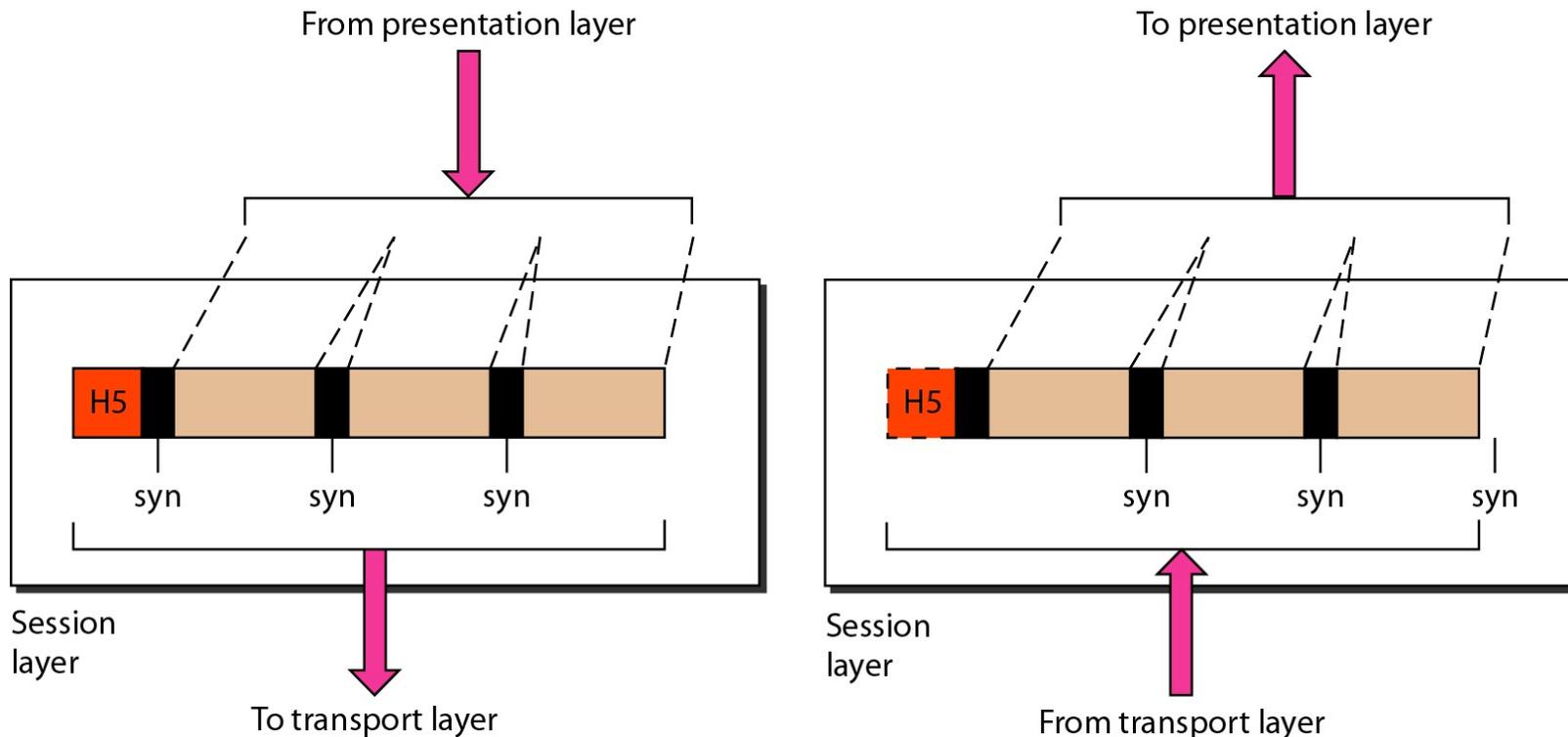


The transport layer is responsible for the delivery of a message from one process to another.

3. Layers In the OSI Models

5) Session Layer

- the session layer is responsible for dialog control and synchronization.
- Dialog control : the Session layer allows two systems to enter into a dialog. It allows the communication between two processes to take place in either half-duplex
- Synchronization : the Session layer allows a process to add checkpoints, or synchronization points, to a stream of data.



3. Layers In the OSI Models

5) Session Layer

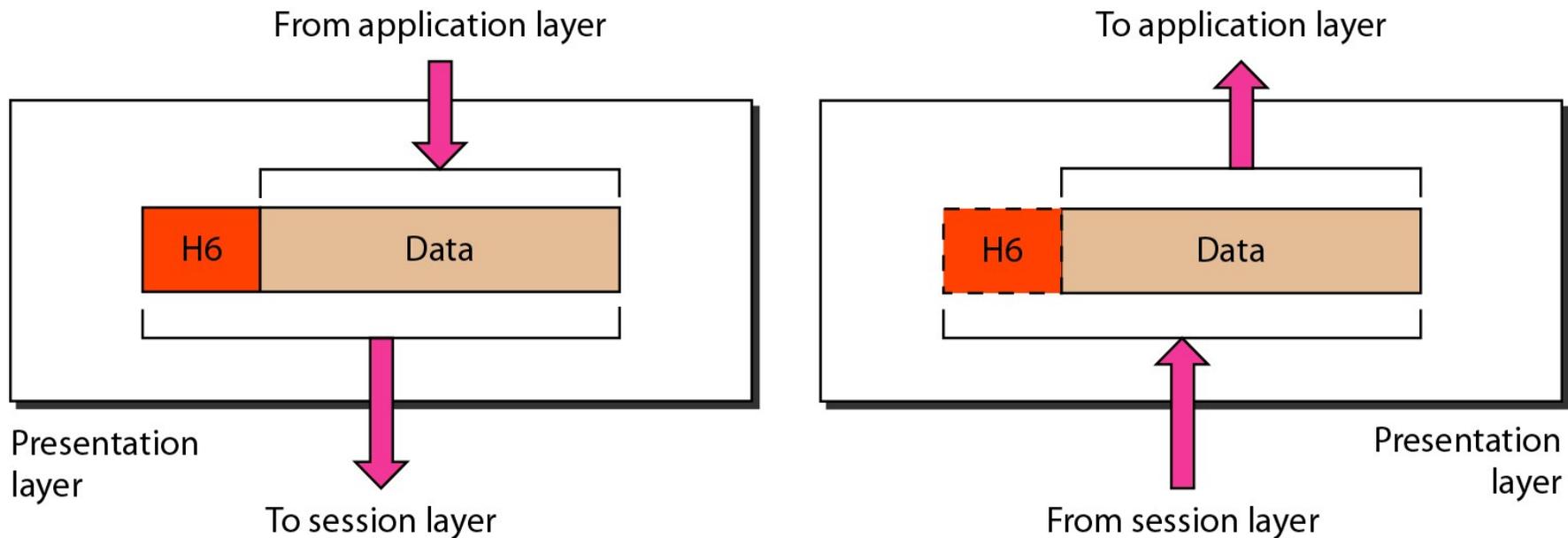
Note

The session layer is responsible for dialog control and synchronization.

3. Layers In the OSI Models

6) Presentation Layer

- responsible for translation, compression, and encryption.
- Translation : the processes in two systems are usually exchanging information in the form of character string, number, and so on.
- Encryption : to carry sensitive information, a system must be able to ensure privacy.
- Compression : Data compression reduces the number of bits contained in the information



3. Layers In the OSI Models

6) Presentation Layer

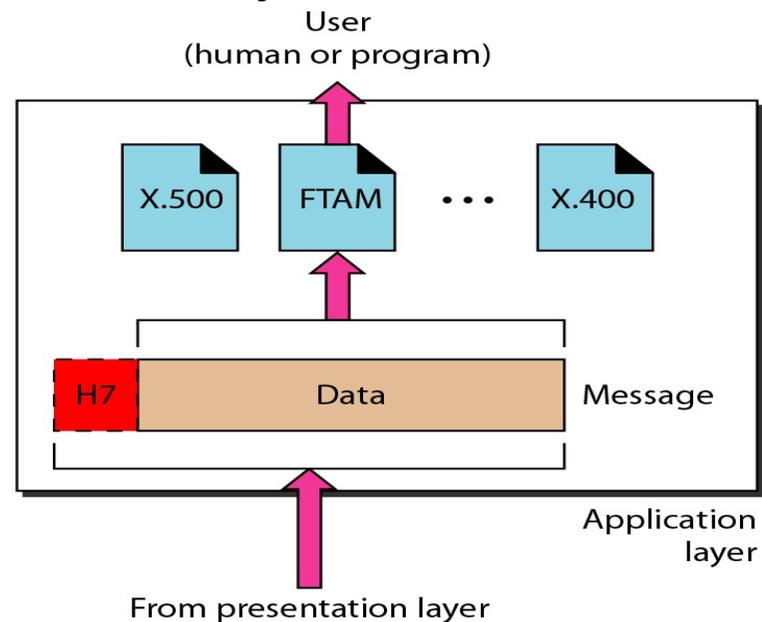
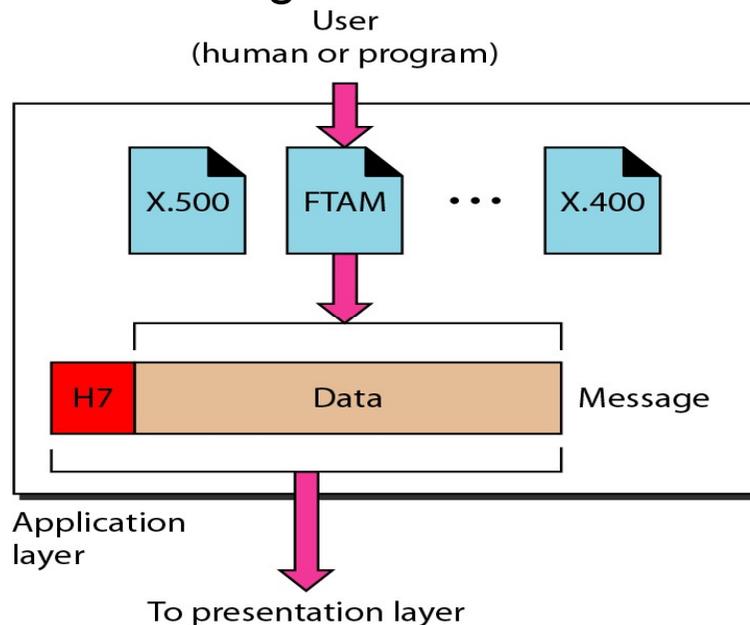


The presentation layer is responsible for translation, compression, and encryption.

3. Layers In the OSI Models

7) Application Layer

- responsible for providing services to the user.
- **Network virtual terminal** : a network virtual terminal is a software version of a physical terminal, and it allows a user to log on to a remote host.
- **File transfer, access, and management** : this application allows a user to access files in a remote host, to retrieve files from a remote computer for use in the local computer, and to manage or control files in a remote computer locally
- **Mail services** : this application provides distributed database sources and access for global information about various objects and services.



3. Layers In the OSI Models

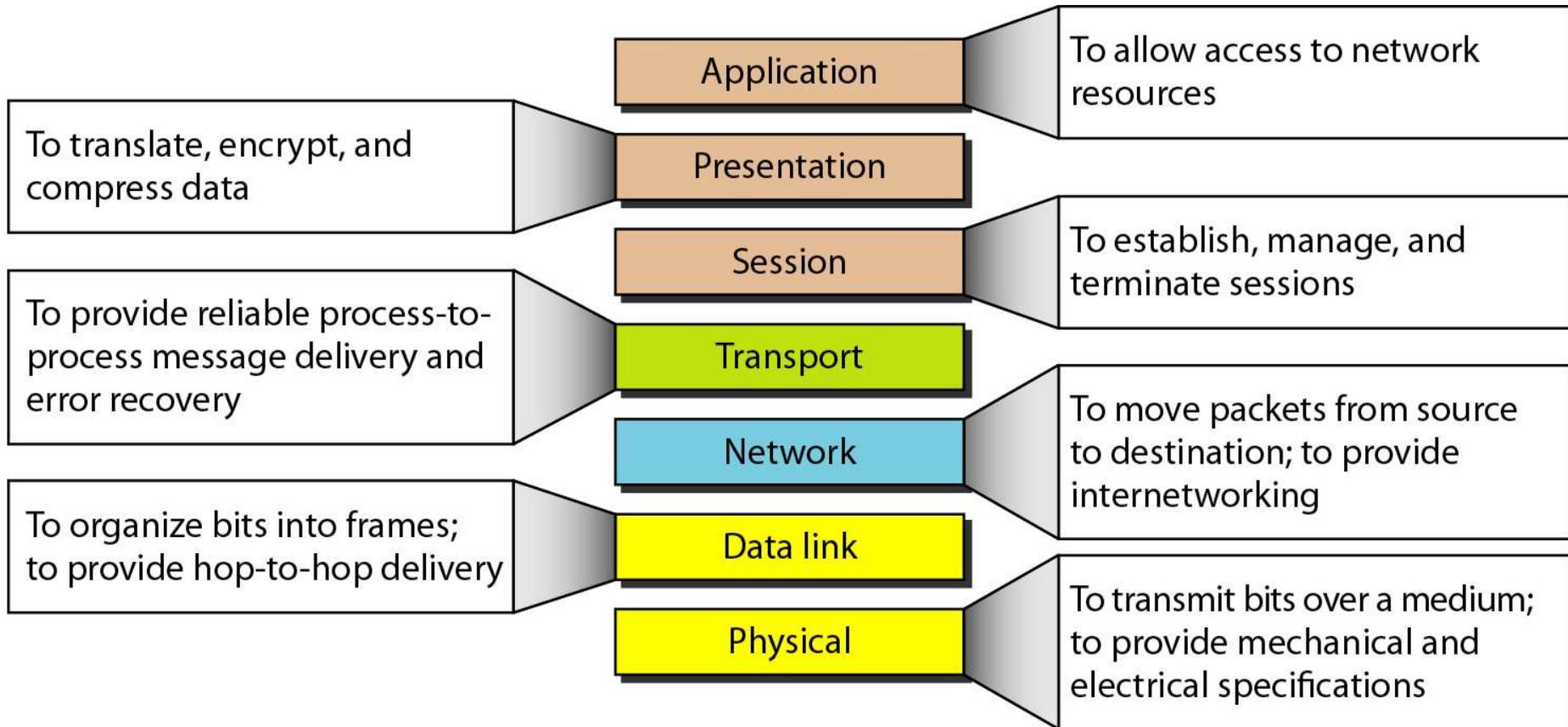
7) Application Layer



The application layer is responsible for providing services to the user.

3. Layers In the OSI Models

Summary of Layers



4. TCP/IP Protocol Suite

*The layers in the **TCP/IP protocol suite** do not exactly match those in the OSI model. The original TCP/IP protocol suite was defined as having four layers: **host-to-network, internet, transport, and application**. However, when TCP/IP is compared to OSI, we can say that the TCP/IP protocol suite is made of five layers: **physical, data link, network, transport, and application**.*

Topics discussed in this section:

Physical and Data Link Layers

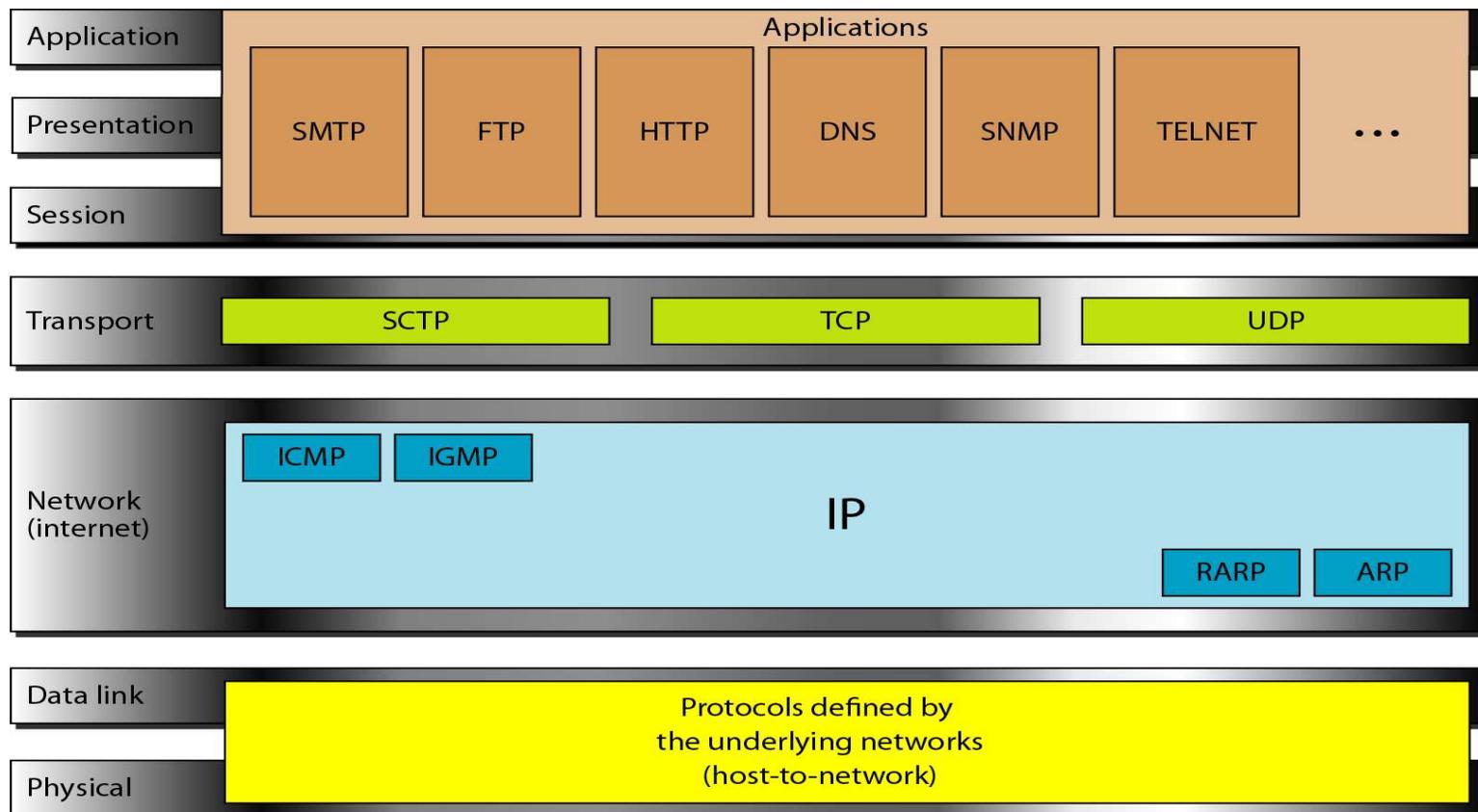
Network Layer

Transport Layer

Application Layer

4. TCP/IP Protocol Suite

The layers in the TCP/IP protocol suite do not exactly match those in the OSI model. The original TCP/IP protocol suite was defined as having four layers: host-to-network, internet, transport, and application. However, when TCP/IP is compared to OSI, we can say that the TCP/IP protocol suite is made of five layers: physical, data link, network, transport, and application.



4. TCP/IP Protocol Suite

■ Physical and Data Link Layers

- At the physical and data link layer, TCP/IP does not define any specific protocol. It supports all the standard and proprietary protocols.

■ Network layer

- Supports the Internetworking Protocol. IP, in turn, uses four supporting protocols, ARP, RARP, ICMP, and IGMP.
- Internetworking Protocol (IP)
 - the transmission mechanism used by the TCP/IP protocols.
 - an unreliable and connectionless protocol - a best-effort delivery service.
 - IP transports data in packets called datagrams
 - The limited functionality of IP should not be considered a weakness, however IP provides bare-bones transmission functions that free the user to add only those facilities necessary for a given application and thereby allows for maximum efficiency.
- Address Resolution Protocol (ARP)
 - ARP is used to associate a logical address with a physical address
 - ARP is used to find the physical address of the node when its Internet address is known.

4. TCP/IP Protocol Suite

- Reverse Address Resolution Protocol (RARP)
 - RARP allows a host to discover its Internet address when it knows only its physical address.
- Internet Control Message Protocol (ICMP)
 - ICMP is a mechanism used by host and gateways to send notification of datagram problems back to the sender.
- Internet Group Message Protocol (IGMP)
 - IGMP is used to facilitate the simultaneous transmission of a message to group of recipients.

4. TCP/IP Protocol Suite

■ Transport layer

- Represented in TCP/IP by two protocols : TCP and UDP
- IP is a host-host protocol, it can deliver a packet from one physical device to another.
- TCP and UDP are transport level protocols responsible for delivery of a message from a process to another process.
- User Datagram Protocol (UDP)
 - UDP is the simpler of the two standard TCP/IP transport protocols
 - A process-to-process protocol that adds only port addresses, checksum error control, and length information to the data from the upper layer.
- Transmission Control Protocol (TCP)
 - TCP provides full transport-layer services to applications.
 - TCP is a reliable stream transport protocol. Cf. The term stream means a connection-oriented.
- Stream Control Transmission Protocol (SCTP)
 - SCTP provides support for newer applications such as voice over the Internet protocol (VoIP).

4. TCP/IP Protocol Suite

■ Application Layer

- The application layer in TCP/IP is equivalent to the combined session, presentation, and application layers in the OSI model.
- Many protocols are defined at this layer.

Summary

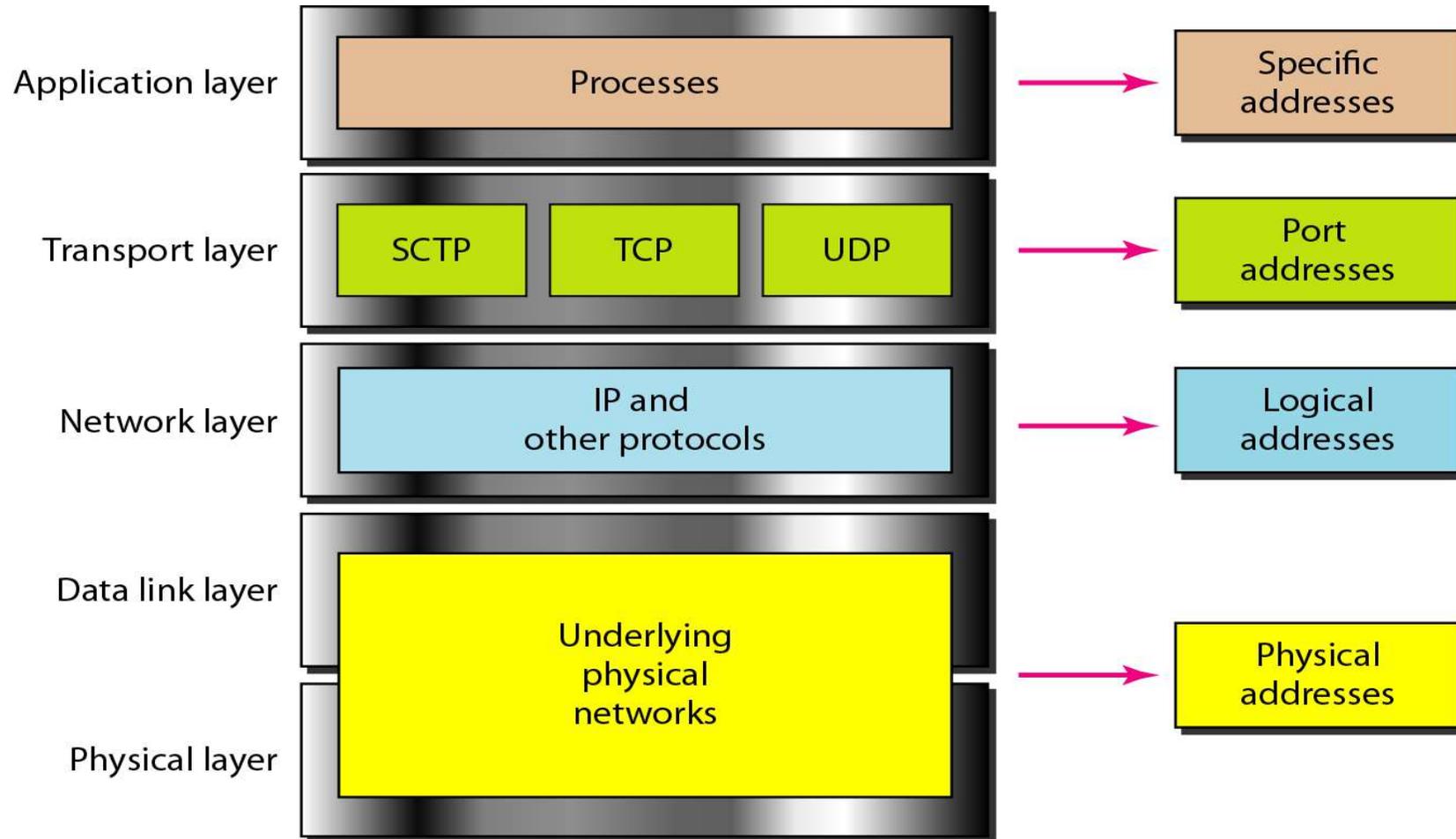
- The OSI-7 model defines seven layers: physical layer, data link layer, network layer, transport layer, session layer, presentation layer, and application layer.
- The physical layer coordinates the functions required to transmit a bit stream over a physical medium.
- The data link layer is responsible for delivering data units from one station to the next without errors.
- The network layer is responsible for the source-to-destination delivery of a packet across multiple network links.
- The transport layer is responsible for the process-to-process delivery of the entire message.
- The application layer enables the users to access the network.

Appendix

- TCP/IP model -



1) Addresses in TCP/IP Model



2) Ethernet

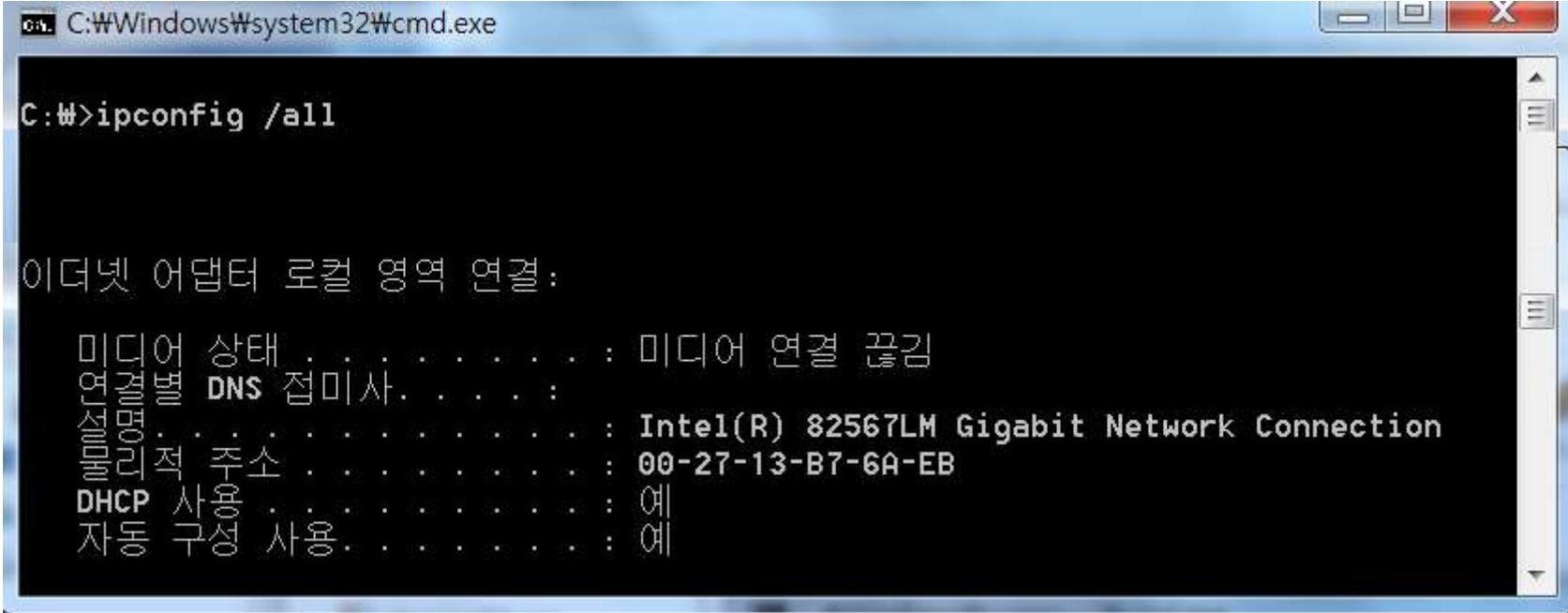
- A frame-based networking technologies for local area networks (LAN).
- Standardization of LAN: IEEE 802.3
- Speed: 10Mbps → 100Mbps → 1Gbps → 10Gbps → 40Gbps/100Gbps
- IEEE 802.3 Ethernet frame structure

Preamble	Start of frame delimiter	MAC destination	MAC source	802.1Q tag (optional)	Ethertype or length	Payload	Cyclic redundancy check	Interframe gap
7 octets of 10101010	1 octet of 10101011	6 octets	6 octets	(4 octets)	2 octets	46–1500 octets	4 octets	12 octets
		64–1522 octets						
		72–1530 octets						
		84–1542 octets						

Physical Address

- Physical Address = Ethernet Address
 - The address of a node as defined by its LAN.
 - 48-bit (6-byte) physical address written as 12 hexadecimal digits; as shown below:

00-27-13-B7-6A-EB



```
C:\Windows\system32\cmd.exe

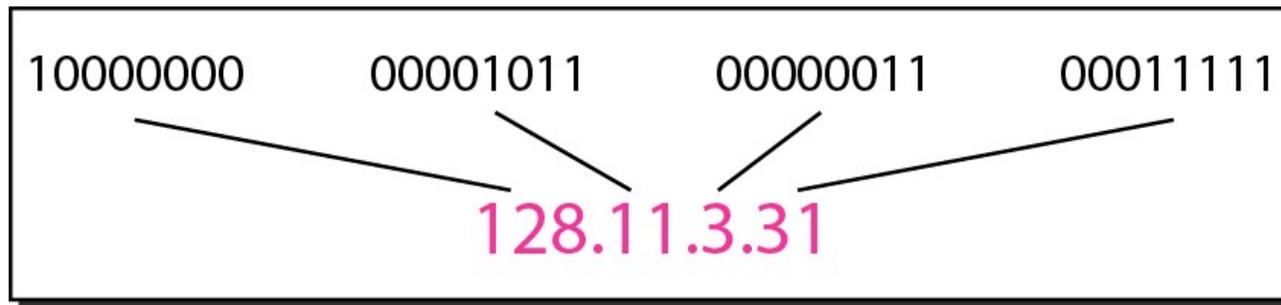
C:\#>ipconfig /all

이더넷 어댑터 로컬 영역 연결:

    미디어 상태 . . . . . : 미디어 연결 끊김
    연결별 DNS 접미사 . . . . :
    설명 . . . . . : Intel(R) 82567LM Gigabit Network Connection
    물리적 주소 . . . . . : 00-27-13-B7-6A-EB
    DHCP 사용 . . . . . : 예
    자동 구성 사용 . . . . . : 예
```

3) IP (Internet Protocol)

- The transmission mechanism used by the TCP/IP protocols.
- An unreliable and connectionless protocol - a best-effort delivery service.
- IP transports data in packets called datagrams.
- The IPv4 addresses are unique and universal.
 - IPv4: 32 bits (4 byte), 2^{32} or 4,294,967,296.
 - IPv6: 128 bits (16 byte): 2^{128}



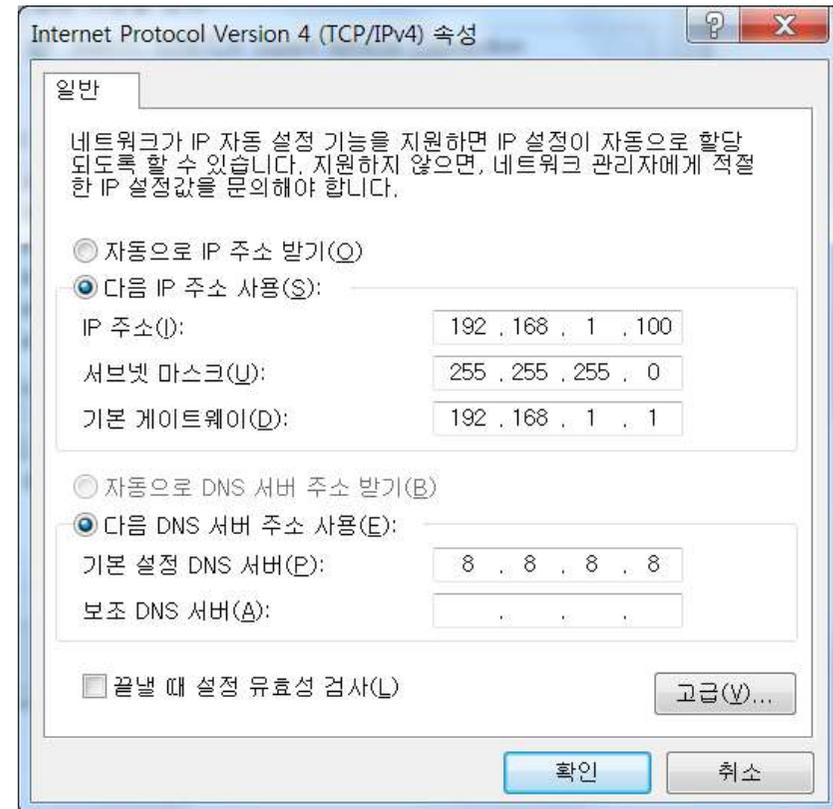
IP Address

- The Internet Assigned Numbers Authority (IANA) manages the IP address space allocations.

	First byte	Second byte	Third byte	Fourth byte	
Class A	0–127				– Class A: large organizations.
Class B	128–191				– Class B: midsize organizations.
Class C	192–223				– Class C: small organizations.
Class D	224–239				– Class D: multicasting.
Class E	240–255				– Class E: reserved for future use.

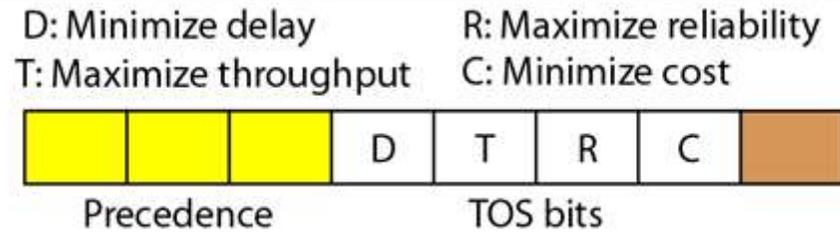
Subnet Mask

- Network part: '1' of subnet mask
- Host part: '0' of subnet mask

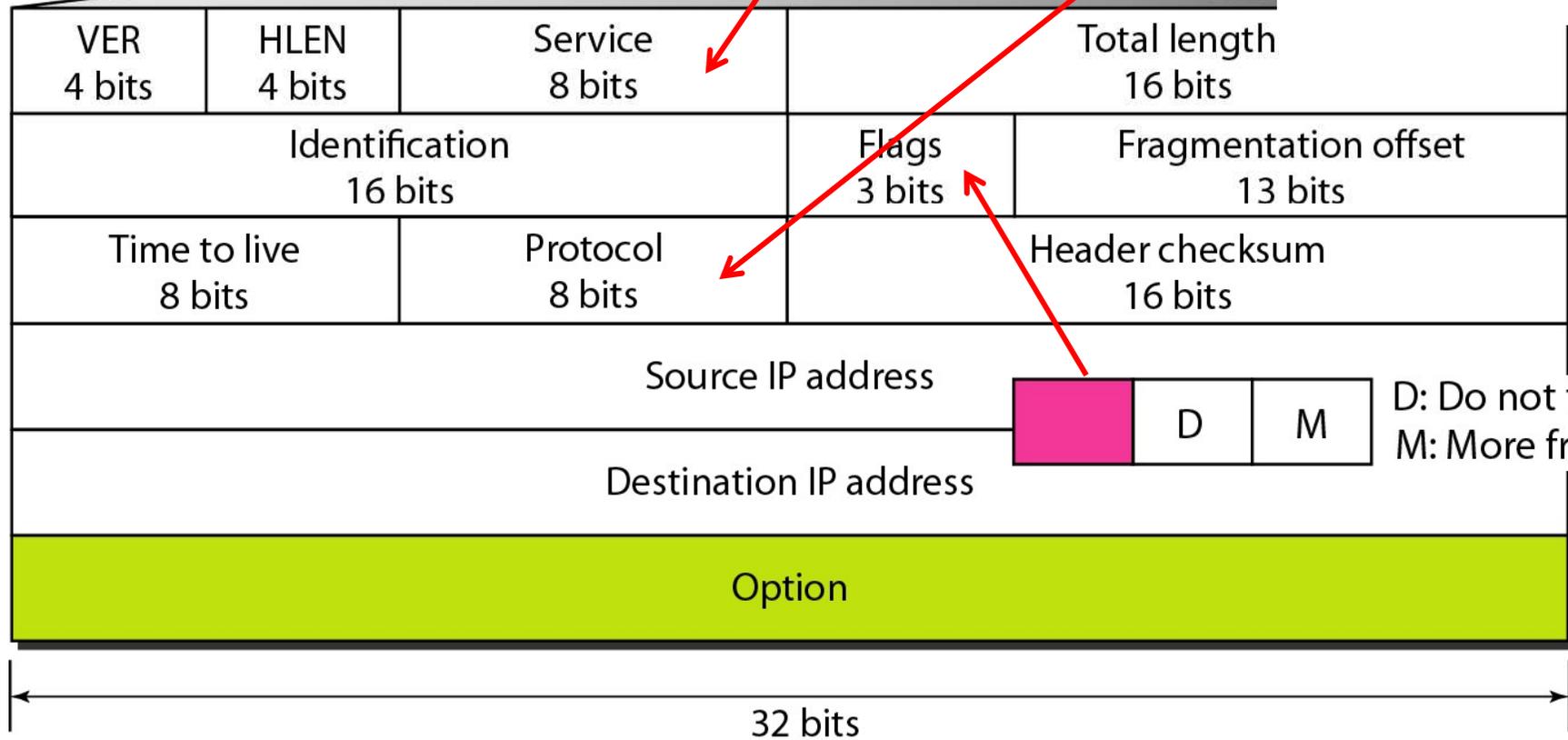
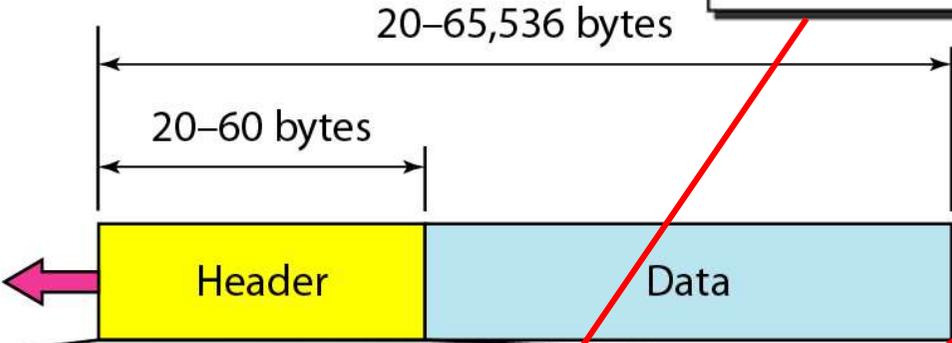


<i>Class</i>	<i>Binary</i>	<i>Dotted-Decimal</i>	<i>CIDR</i>
A	11111111 00000000 00000000 00000000	255.0.0.0	/8
B	11111111 11111111 00000000 00000000	255.255.0.0	/16
C	11111111 11111111 11111111 00000000	255.255.255.0	/24

IP header

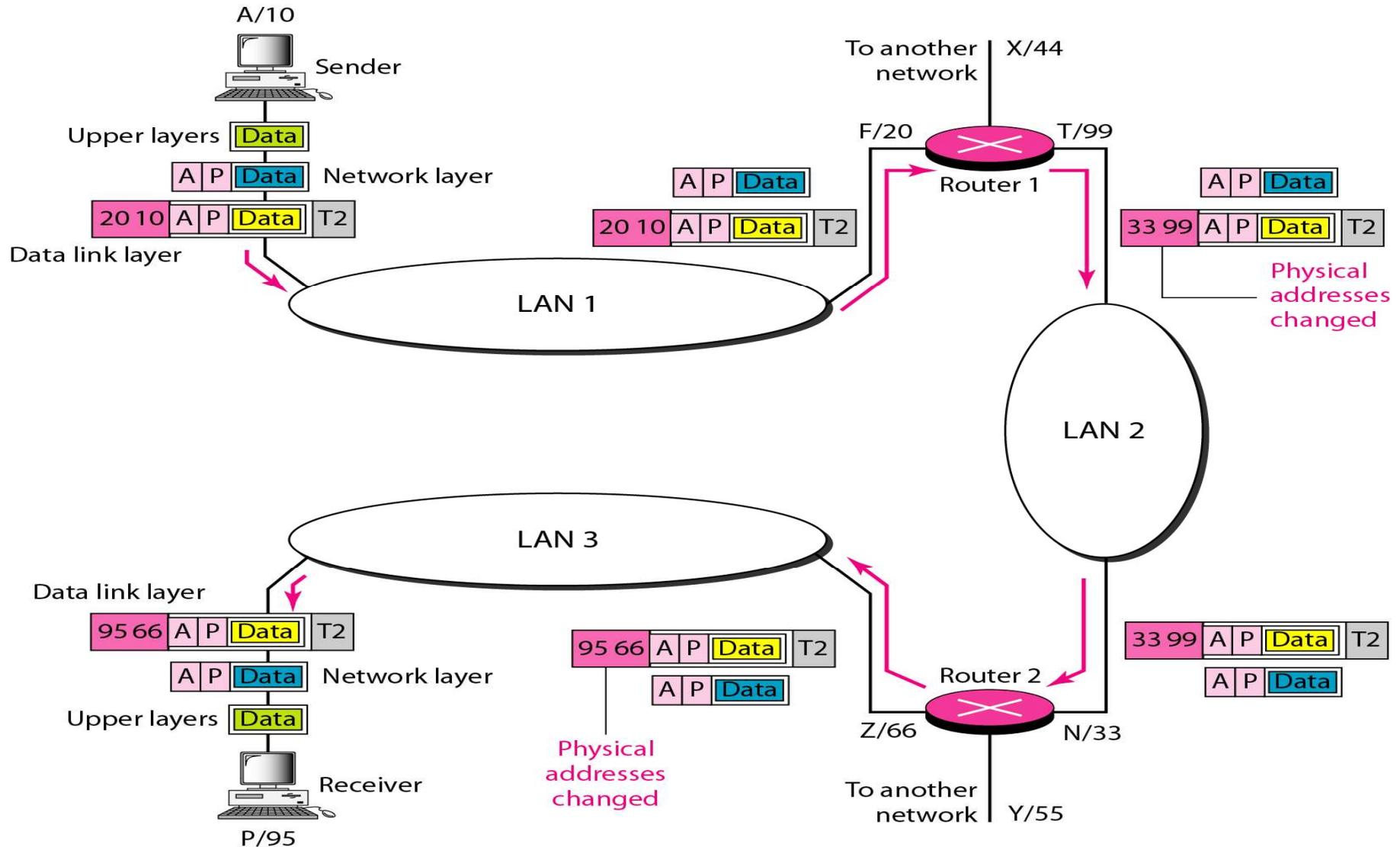


Value	Protocol
1	ICMP
2	IGMP
6	TCP
17	UDP
89	OSPF



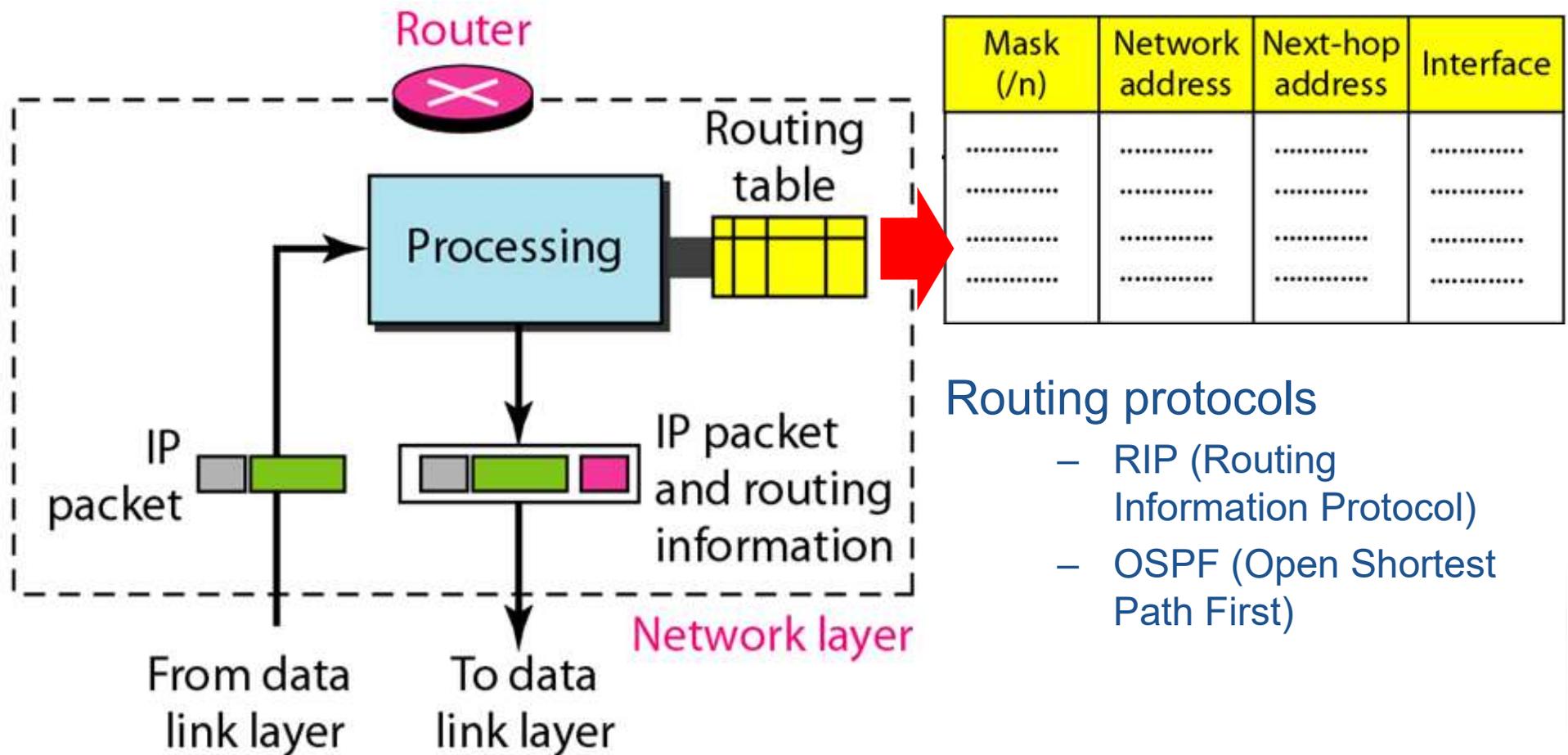
Packet Delivery (1/2)

- The physical addresses will change from hop to hop, but the logical addresses usually remain the same.



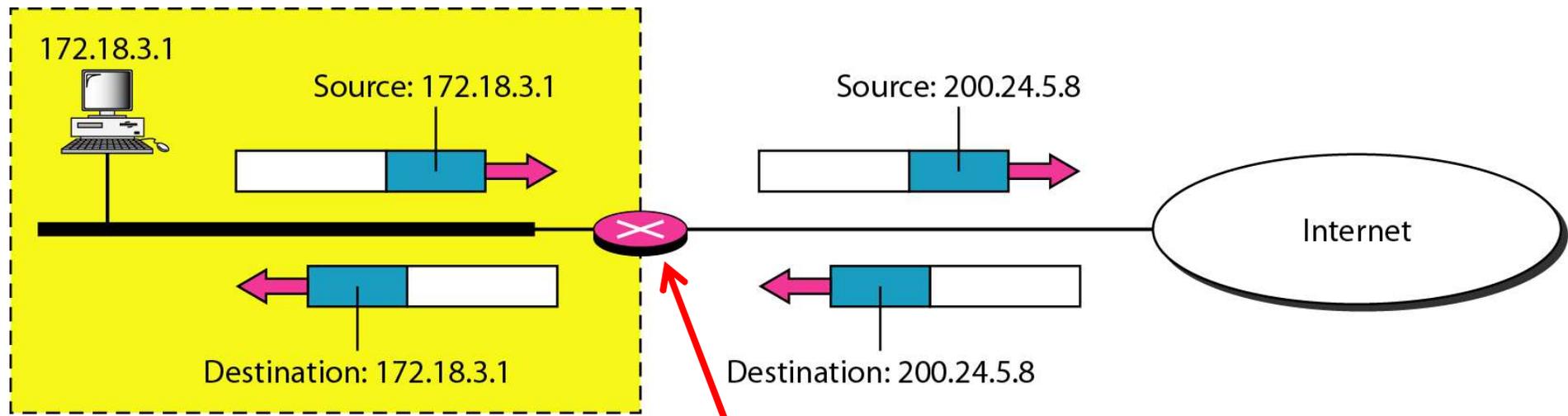
Packet Delivery (2/2)

- Router at the Network layer



NAT (Network Address Translation)

- NAT enables a user to have a large set of addresses internally and one address, or a small set of addresses, externally.

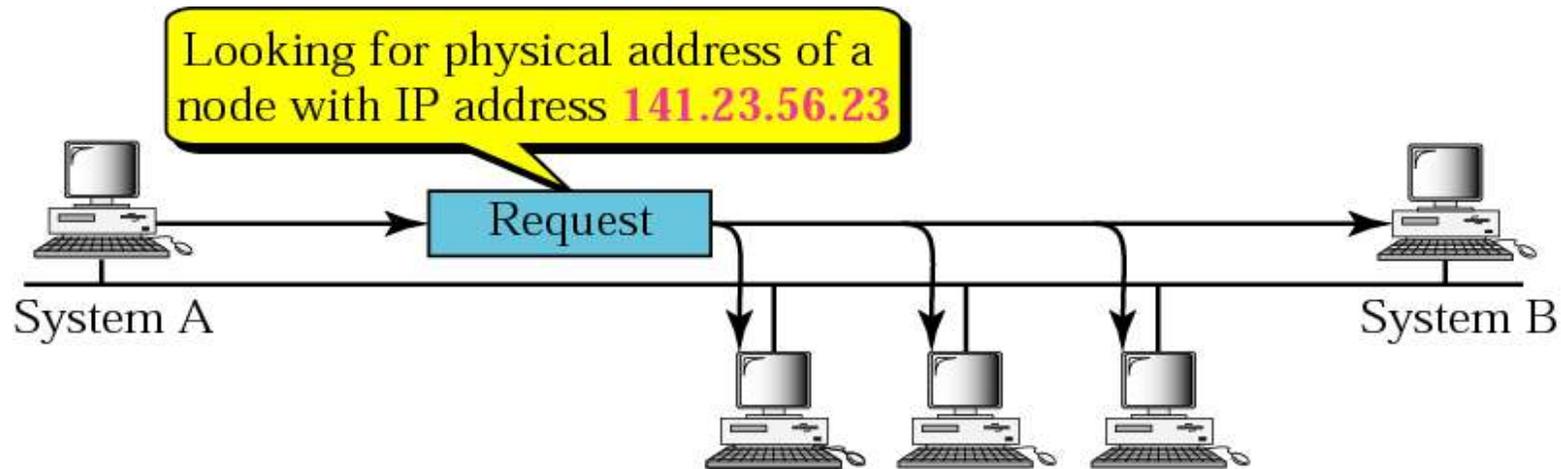


<i>Private Address</i>	<i>Private Port</i>	<i>External Address</i>	<i>External Port</i>	<i>Transport Protocol</i>
172.18.3.1	1400	25.8.3.2	80	TCP
172.18.3.2	1401	25.8.3.2	80	TCP
...

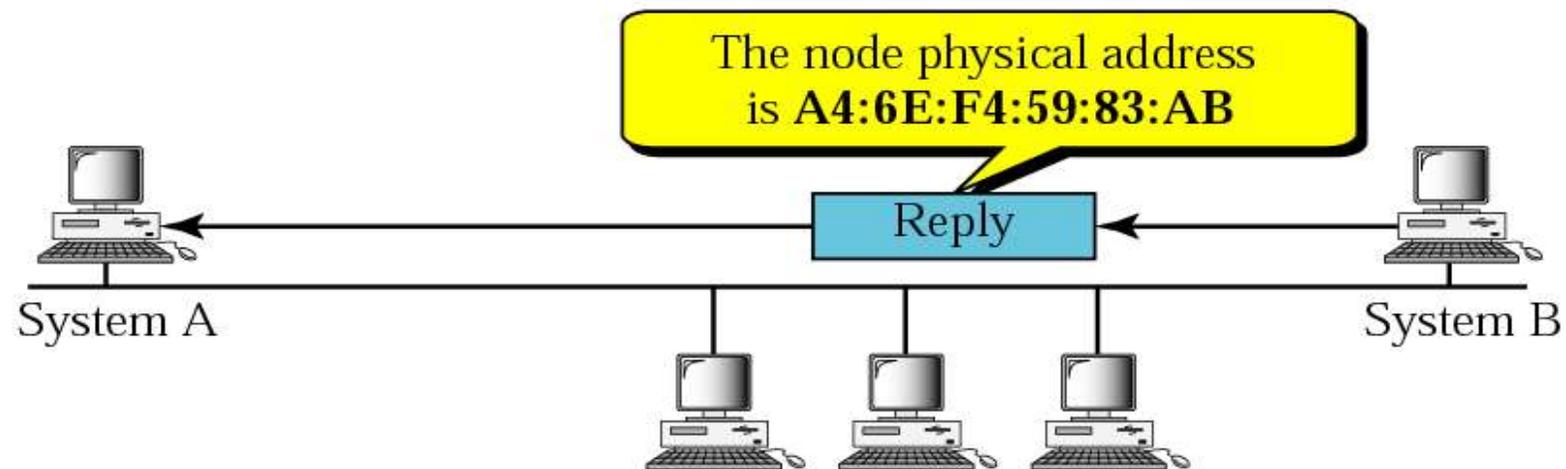
ARP, RARP, ICMP, IGMP

- Address Resolution Protocol (ARP)
 - ARP is used to find the **physical address** of the node when its Internet address is known.
- Reverse Address Resolution Protocol (RARP)
 - RARP allows a host to discover its **Internet address** when it knows only its physical address.
- Internet Control Message Protocol (ICMP)
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ARP (1/2)



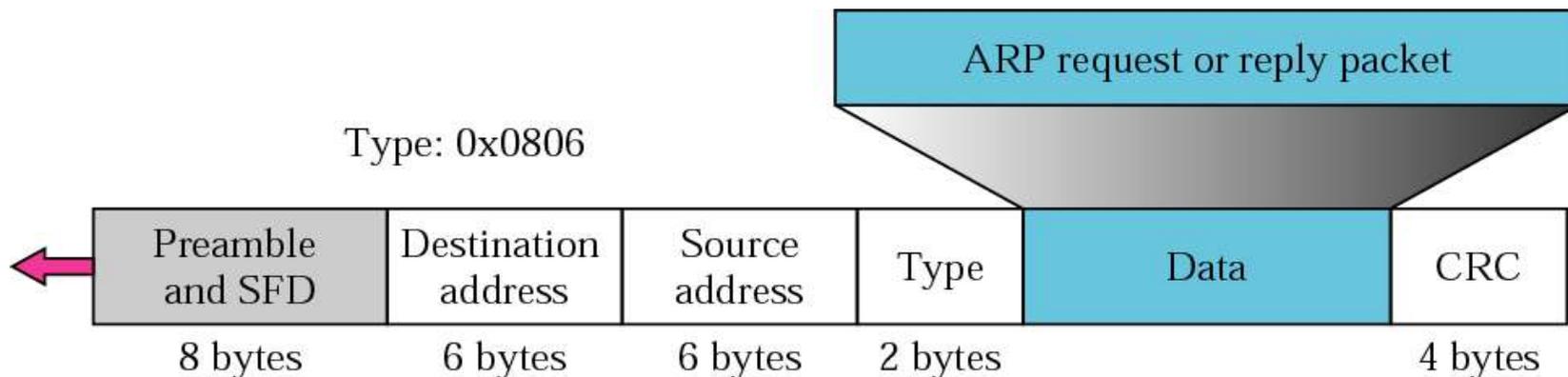
a. ARP request is broadcast



b. ARP reply is unicast

ARP (2/2)

Hardware Type Ethernet (1)		Protocol Type IPv4 (0800)
Hardware length	Protocol length	Operation Request 1, Reply 2
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled in a request)		
Target protocol address (For example, 4 bytes for IP)		



ARP request

The image shows a Wireshark capture window titled "(Untitled) - Wireshark". The interface includes a menu bar (File, Edit, View, Go, Capture, Analyze, Statistics, Help), a toolbar with various icons, and a filter field. The main display area shows a list of captured packets. Packet 27 is selected, showing an ARP request. The details pane below the packet list shows the structure of the Ethernet II frame and the ARP request, including source and target MAC and IP addresses. The packet bytes pane at the bottom shows the raw hex and ASCII data of the selected packet.

No.	Time	Source	Destination	Protocol	Info
27	21.500953	IntelCor_8b:27:07	Broadcast	ARP	who has 192.168.1.1? Tell 192.168.1.210
28	21.502170	EfmNetwo_9a:46:35	IntelCor_8b:27:07	ARP	192.168.1.1 is at 00:08:9f:9a:46:35
29	21.502185	192.168.1.210	192.168.1.1	ICMP	Echo (ping) request
30	21.505167	192.168.1.1	192.168.1.210	ICMP	Echo (ping) reply

Frame 27 (42 bytes on wire, 42 bytes captured)

- Ethernet II, src: IntelCor_8b:27:07 (00:1c:c0:8b:27:07), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
 - Destination: Broadcast (ff:ff:ff:ff:ff:ff)
 - Source: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Type: ARP (0x0806)
- Address Resolution Protocol (request)
 - Hardware type: Ethernet (0x0001)
 - Protocol type: IP (0x0800)
 - Hardware size: 6
 - Protocol size: 4
 - Opcode: request (0x0001)
 - Sender MAC address: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Sender IP address: 192.168.1.210 (192.168.1.210)
 - Target MAC address: 00:00:00_00:00:00 (00:00:00:00:00:00)
 - Target IP address: 192.168.1.1 (192.168.1.1)

```
0000 ff ff ff ff ff 00 1c c0 8b 27 07 08 06 00 01 .....  
0010 08 00 06 04 00 01 00 1c c0 8b 27 07 c0 a8 01 d2 .....  
0020 00 00 00 00 00 00 c0 a8 01 01 .....
```

Ethernet (eth), 14 bytes Packets: 46 Displayed: 46 Marked: 0 Dropped: 0 Profile: Default

ARP reply

The image shows a Wireshark capture window titled "(Untitled) - Wireshark". The interface includes a menu bar (File, Edit, View, Go, Capture, Analyze, Statistics, Help), a toolbar with various icons, and a filter field. The main display area is divided into three sections: a packet list, a packet details pane, and a packet bytes pane.

Packet List:

No.	Time	Source	Destination	Protocol	Info
27	21.500953	IntelCor_8b:27:07	Broadcast	ARP	who has 192.168.1.1? Tell 192.168.1.210
28	21.502170	EfmNetwo_9a:46:35	IntelCor_8b:27:07	ARP	192.168.1.1 is at 00:08:9f:9a:46:35
29	21.502185	192.168.1.210	192.168.1.1	ICMP	Echo (ping) request
30	21.505167	192.168.1.1	192.168.1.210	ICMP	Echo (ping) reply

Packet Details (Frame 28):

- Frame 28 (60 bytes on wire, 60 bytes captured)
- Ethernet II, Src: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35), Dst: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Destination: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Source: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35)
 - Type: ARP (0x0806)
 - Trailer: 00000000000000000000000000000000
- Address Resolution Protocol (reply)
 - Hardware type: Ethernet (0x0001)
 - Protocol type: IP (0x0800)
 - Hardware size: 6
 - Protocol size: 4
 - opcode: reply (0x0002)
 - Sender MAC address: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35)
 - Sender IP address: 192.168.1.1 (192.168.1.1)
 - Target MAC address: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Target IP address: 192.168.1.210 (192.168.1.210)

Packet Bytes:

0000	00 1c c0 8b 27 07	00 08 9f 9a 46 35	08 06 00 01F5....
0010	08 00 06 04 00 02	00 08 9f 9a 46 35	c0 a8 01 01F5....
0020	00 1c c0 8b 27 07	c0 a8 01 d2 00 00	00 00 00 00
0030	00 00 00 00 00 00	00 00 00 00 00 00	

Destination Hardware Address (eth.dst), 6 ... Packets: 46 Displayed: 46 Marked: 0 Dropped: 0 Profile: Default

ICMP (ping request)

The image shows a Wireshark capture window titled "(Untitled) - Wireshark". The main display area shows a list of network packets. Packet 29 is selected, showing an ICMP Echo (ping) request from 192.168.1.210 to 192.168.1.1. Packet 30 is the corresponding ICMP Echo (ping) reply from 192.168.1.1 to 192.168.1.210. The packet details pane for packet 29 shows the following information:

- Frame 29 (74 bytes on wire, 74 bytes captured)
- Ethernet II, Src: IntelCor_8b:27:07 (00:1c:c0:8b:27:07), Dst: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35)
 - Destination: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35)
 - Source: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Type: IP (0x0800)
- Internet Protocol, Src: 192.168.1.210 (192.168.1.210), Dst: 192.168.1.1 (192.168.1.1)
 - Version: 4
 - Header length: 20 bytes
 - Differentiated Services Field: 0x00 (DSCP 0x00: Default; ECN: 0x00)
 - Total Length: 60
 - Identification: 0x51de (20958)
 - Flags: 0x00
 - Fragment offset: 0
 - Time to live: 128
 - Protocol: ICMP (0x01)

The packet bytes pane shows the raw data for the selected packet, with a hex dump on the left and an ASCII representation on the right. The ASCII representation shows the ICMP Echo request structure: `....F5.. ..'...E.
<Q.....
..MZ.. ..abcdef
ghijklmn opqrstuv
wabcdefg hi`

At the bottom of the window, the status bar indicates: "Internet Protocol (ip), 20 bytes" | "Packets: 46 Displayed: 46 Marked: 0 Dropped: 0" | "Profile: Default".

ICMP (ping reply)

The image shows a Wireshark capture of an ICMP Echo (ping) reply. The packet list shows four packets: an ARP request for 192.168.1.1, an ARP response, an ICMP Echo request, and an ICMP Echo reply. The selected packet (No. 30) is an ICMP Echo reply from 192.168.1.1 to 192.168.1.210. The packet details show it is an Internet Protocol (IP) packet with a version of 4, header length of 20 bytes, and protocol of ICMP (0x01). The packet bytes pane shows the raw data in hexadecimal and ASCII.

No.	Time	Source	Destination	Protocol	Info
27	21.500953	IntelCor_8b:27:07	Broadcast	ARP	who has 192.168.1.1? Tell 192.168.1.210
28	21.502170	EfmNetwo_9a:46:35	IntelCor_8b:27:07	ARP	192.168.1.1 is at 00:08:9f:9a:46:35
29	21.502185	192.168.1.210	192.168.1.1	ICMP	Echo (ping) request
30	21.505167	192.168.1.1	192.168.1.210	ICMP	Echo (ping) reply

Frame 30 (74 bytes on wire, 74 bytes captured)

- Ethernet II, Src: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35), Dst: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Destination: IntelCor_8b:27:07 (00:1c:c0:8b:27:07)
 - Source: EfmNetwo_9a:46:35 (00:08:9f:9a:46:35)
 - Type: IP (0x0800)
- Internet Protocol, Src: 192.168.1.1 (192.168.1.1), Dst: 192.168.1.210 (192.168.1.210)
 - Version: 4
 - Header length: 20 bytes
 - Differentiated services Field: 0x00 (DSCP 0x00: Default; ECN: 0x00)
 - Total Length: 60
 - Identification: 0xb79b (47003)
 - Flags: 0x00
 - Fragment offset: 0
 - Time to live: 64
 - Protocol: ICMP (0x01)

```
0000  00 1c c0 8b 27 07 00 08 9f 9a 46 35 08 00 45 00  ....'... ..F5..E.
0010  00 3c b7 9b 00 00 40 01 3f 02 c0 a8 01 01 c0 a8  .<....@. ?.....
0020  01 d2 00 00 55 5a 00 01 00 01 61 62 63 64 65 66  ..UZ.. ..abcdef
0030  67 68 69 6a 6b 6c 6d 6e 6f 70 71 72 73 74 75 76  ghijklmn opqrstuv
0040  77 61 62 63 64 65 66 67 68 69  wabcdefg hi
```

Internet Protocol (ip), 20 bytes Packets: 46 Displayed: 46 Marked: 0 Dropped: 0 Profile: Default

Quiz

- What is the source IP address of the packet?

```
0000  00 08 9f 9a 46 35 00 24  54 30 e8 1f 08 00 45 00
0010  05 74 02 22 20 00 80 01  00 00 c0 a8 01 05 c0 a8
0020  01 01 08 00 e4 4b 00 01  01 9e 61 62 63 64 65 66
0030  67 68 69 6a 6b 6c 6d 6e  6f 70 71 72 73 74 75 76
0040  77 61 62 63 64 65 66 67  68 69 6a 6b 6c 6d 6e 6f
0050  70 71 72 73 74 75 76 77  61 62 63 64 65 66 67 68
0060  69 6a 6b 6c 6d 6e 6f 70  71 72 73 74 75 76 77 61
0070  62 63 64 65 66 67 68 69  6a 6b 6c 6d 6e 6f 70 71
```

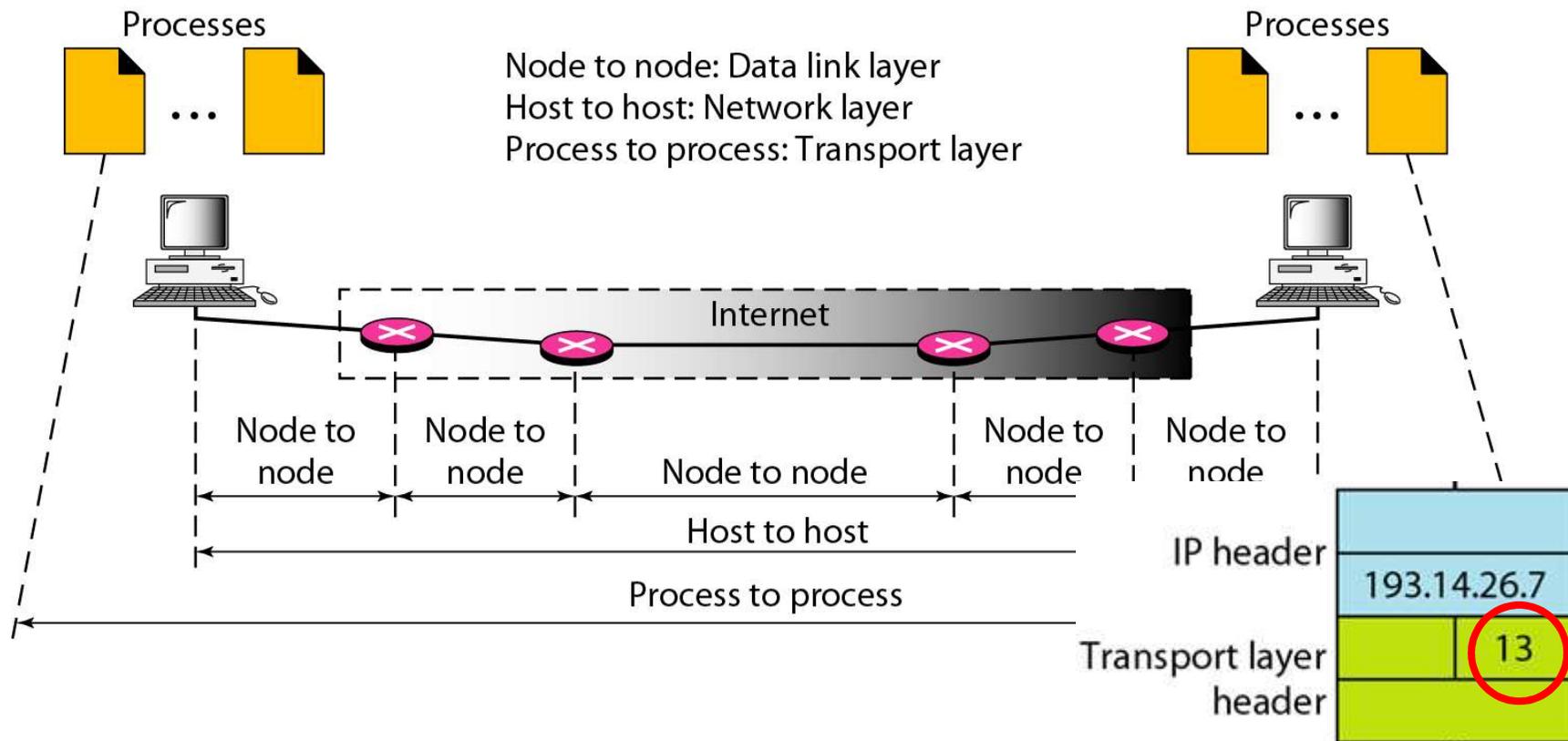
Answer: ?

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address (32 bits)				
Destination IP address (32 bits)				

4) TCP/UDP

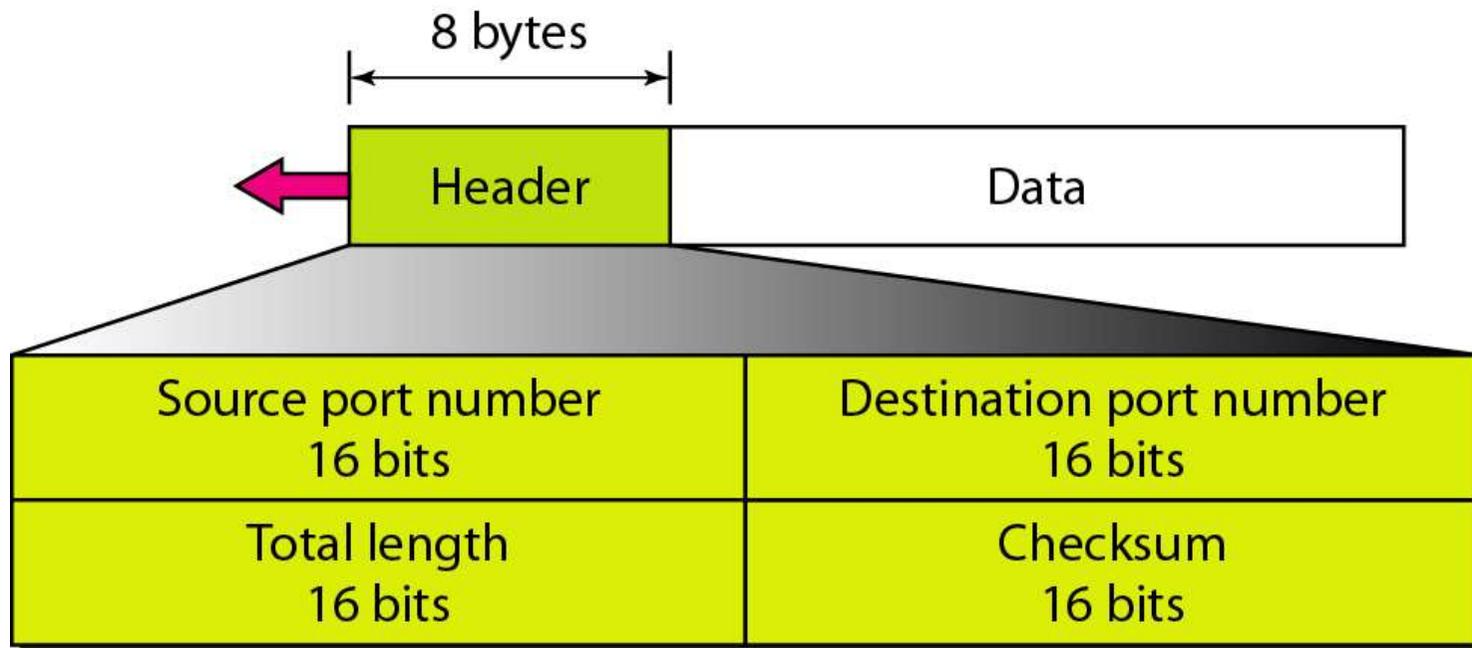
- Responsible for delivery of a message from a process to another process.
- User Datagram Protocol (UDP)
 - UDP is the simpler of the two standard TCP/IP transport protocols
- Transmission Control Protocol (TCP)
 - TCP is a reliable stream transport protocol.
 - Connection-oriented protocol.
- Stream Control Transmission Protocol (SCTP)
 - SCTP provides support for newer applications such as voice over the Internet protocol (VoIP).

Port Address



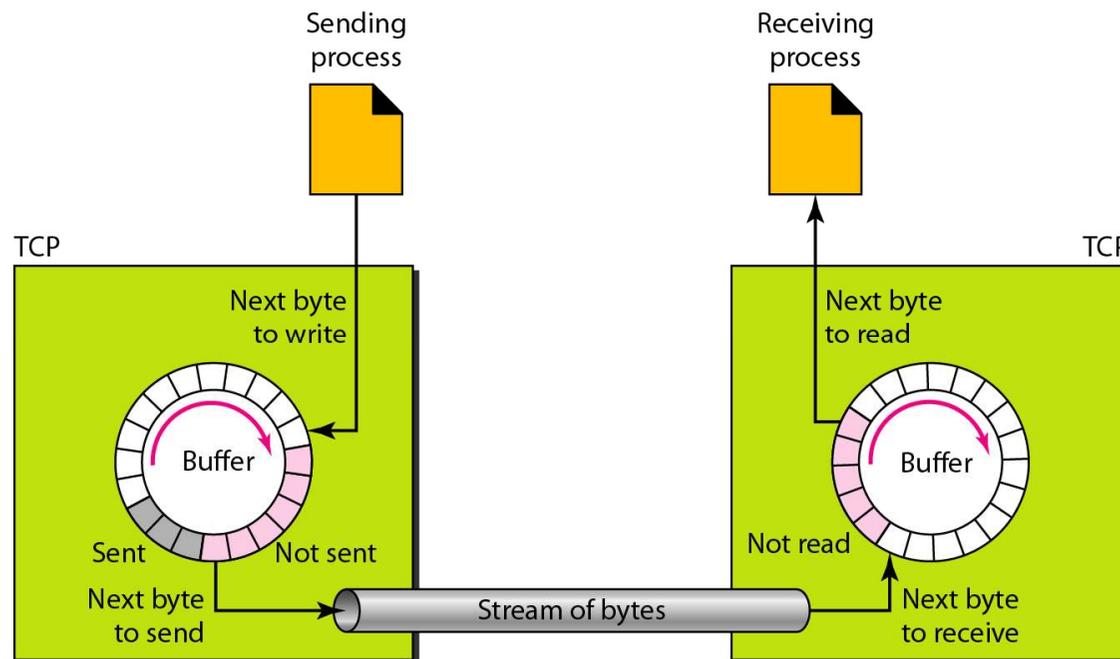
UDP

- The User Datagram Protocol (UDP) is called a connectionless, unreliable transport protocol.
- Protocol: Echo (7), BOOTPs/c (67, 68), TFTP (69), NTP (123), SNMP (161), SNMP trap(162)

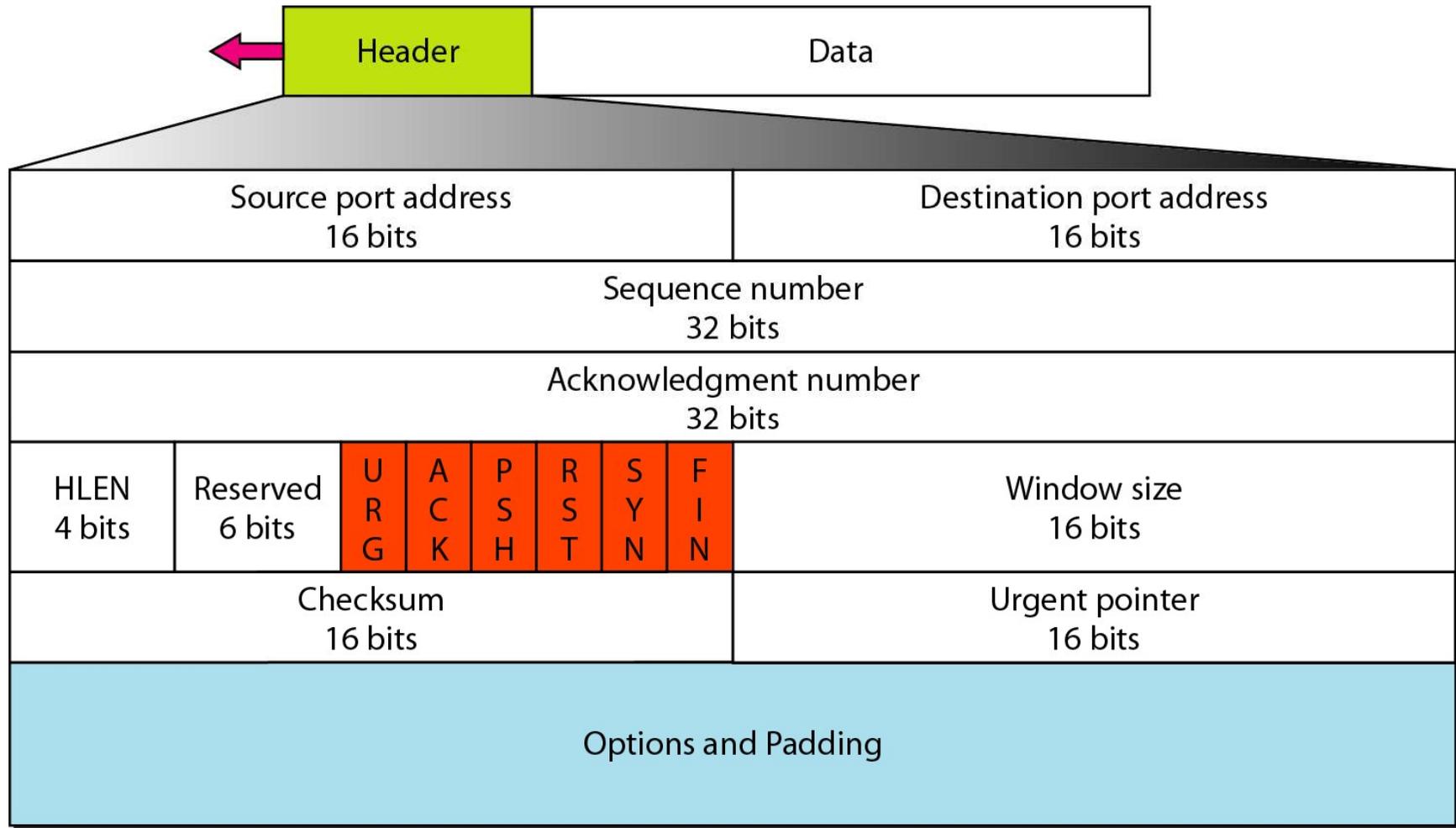


TCP

- Connection-oriented protocol.
- TCP creates a virtual connection between two TCPs to send data.
- Congestion control: Slow start, Reno, Vegas, ...
- Protocol: FTP (21), TELNET (23), SMTP (25), DNS (53), BOOTP (67), HTTP (80)



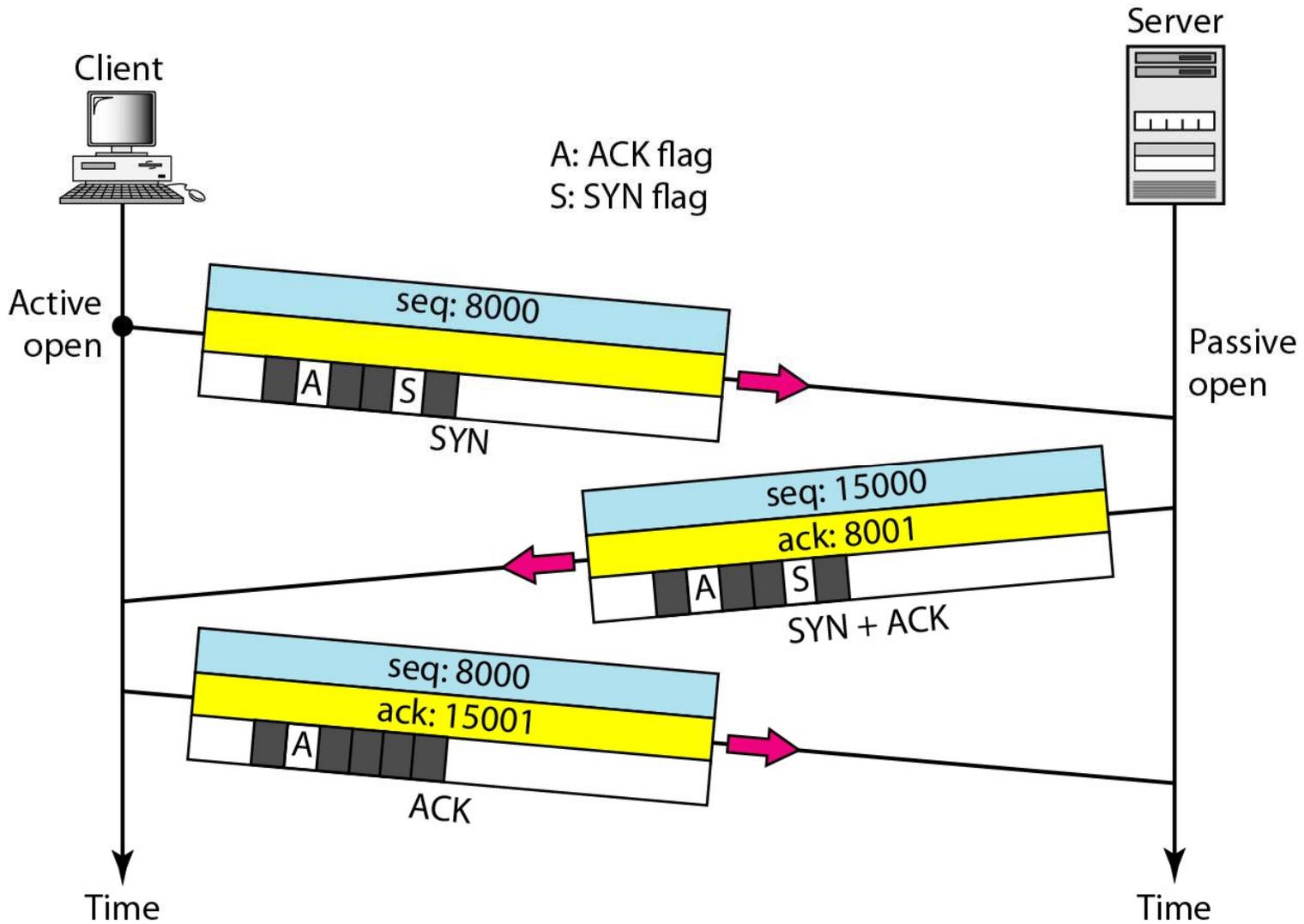
TCP Header



URG: Urgent pointer is valid
 ACK: Acknowledgment is valid
 PSH: Request for push

RST: Reset the connection
 SYN: Synchronize sequence numbers
 FIN: Terminate the connection

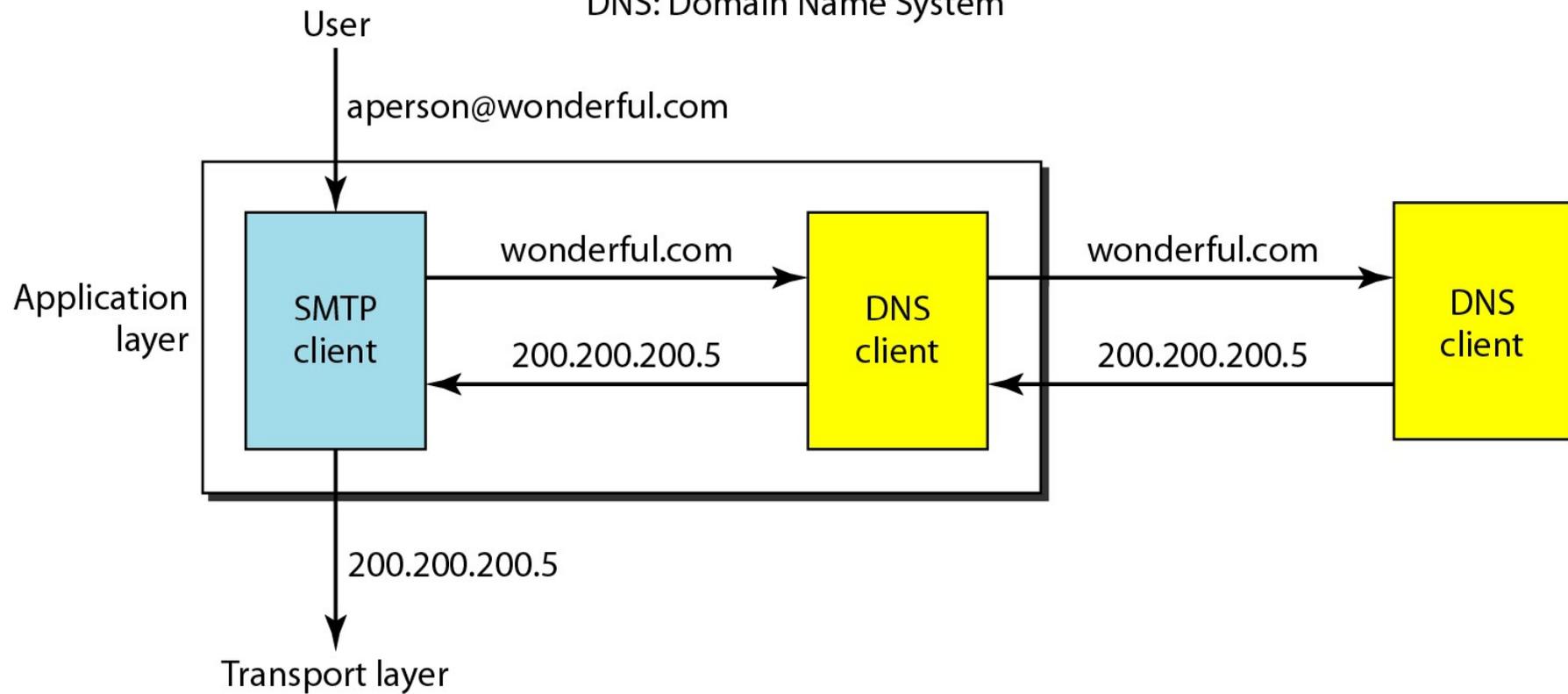
TCP: Connection Establishment



5) DNS, HTTP

- DNS Service

SMTP: Simple Mail Transfer Protocol (e-mail)
DNS: Domain Name System



Architecture of WWW

